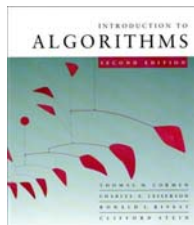




CS 2233 -- Fall 2009



Graphs

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Slides courtesy of Charles Leiserson with changes and additions by Carola Wenk

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1



Graphs

Definition.

A (simple) **undirected graph** $G = (V, E)$ consists of

- a set V of **vertices** (singular: **vertex**), and
- a set E of **edges** consisting of 2-element subsets of V .

unordered pairs

A **directed graph (digraph)** $G = (V, E)$ consists of

- a set V of **vertices**, and
- a set $E \subseteq V \times V$ of **edges**.

ordered pairs

In either case, we have $|E| \in O(|V|^2)$.

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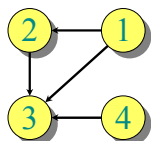
2



Adjacency-matrix representation

The **adjacency matrix** of a graph $G = (V, E)$, where $V = \{1, 2, \dots, n\}$, is the matrix $A[1 \dots n, 1 \dots n]$ given by

$$A[i, j] = \begin{cases} 1 & \text{if } (i, j) \in E, \\ 0 & \text{if } (i, j) \notin E. \end{cases}$$



| A | 1 | 2 | 3 | 4 |
|-----|---|---|---|---|
| 1 | 0 | 1 | 1 | 0 |
| 2 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 1 | 0 |

$\Theta(|V|^2)$ storage
 \Rightarrow **dense**
 representation.

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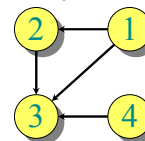
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3



Adjacency-list representation

An **adjacency list** of a vertex $v \in V$ is the list $Adj[v]$ of vertices adjacent to v .



$Adj[1] = \{2, 3\}$
 $Adj[2] = \{3\}$
 $Adj[3] = \{1, 2, 4\}$
 $Adj[4] = \{3\}$

degree

For undirected graphs, $|Adj[v]| = \deg(v)$.

For digraphs, $|Adj[v]| = \deg^+(v)$.

out-degree

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Adjacency-list representation

Handshaking Lemma:

Every edge is counted twice

- For undirected graphs:

$$\sum_{v \in V} \deg(v) = 2|E|$$

- For digraphs:

$$\sum_{v \in V} \deg^-(v) + \sum_{v \in V} \deg^+(v) = 2|E|$$

in-degree

out-degree

⇒ adjacency lists use $\Theta(|V| + |E|)$ storage

⇒ a ***sparse*** representation

⇒ We usually use this representation,
unless stated otherwise



Paths, Cycles, Connectivity

Let $G=(V,E)$ be a directed (or undirected) graph

- A **path** of length k in G is a sequence of vertices v_1, v_2, \dots, v_{k+1} such that $(v_i, v_{i+1}) \in E$ (or $\{v_i, v_{i+1}\} \in E$ if G is undirected) for all $i \in \{1, \dots, k\}$.
- A path is **simple** if all vertices in the path are distinct.
- A path v_1, v_2, \dots, v_{k+1} forms a **cycle** if $v_1 = v_{k+1}$ and $k \geq 2$.
- An undirected graph is **connected** if for every pair $u, v \in V$ there is a path from u to v . (Note that then there is also a path from v to u .)
- A directed graph is **strongly connected** if for every pair $u, v \in V$ there is a path from u to v and there is a path from v to u .
- The **(strongly) connected components** of a graph are the equivalence classes of vertices under this reachability relation.
- A graph with no cycles is **acyclic**.
 - An undirected connected acyclic graph is called a **tree**.
 - An undirected acyclic graph is a **forest** (= set of trees).
 - A **rooted tree** is a tree that has one distinguished **root** vertex.