#### CMPS 6630: Introduction to Computational Biology and Bioinformatics

Multiple Sequence Alignment and Phylogenetic Trees

# Sequence Alignment

- So Far:
  - -Global Alignment
  - -Local Alignment
  - -Database Matching
- All were for a single pair of sequences
- What if we want to look at more than 2 sequences?



# Multiple Sequence Alignment

gi|Tetrahymenathenrophila/1-439 gi|Schizosaccharonycespon/be/1-395 gi|Prochlorococcusmarinus/1-384 gi|toxoplasmagondii/1-536 gi|Plasmodium3D7falciparum/1-527 gi|Plasmodiumberghei/1-506 gi|Plasmodiumberghei/1-506 gi|Plasmodiumchabaudi/1-499 gi|Plasmodiumvivax/1-547

MSTFGRIYRVSTFGESHSPAVGCTVEGFPSLIKVQREEIDKALDRRPGDSSLTTKRNETDDVEILCOLENEISIGSPIAFLVRNDDQIK
MSSFGTLFKVTTYGESHCKSVGCIVEGCPPGMNLTESDVDVDLTRRPGDSNLTTPRNEKDKVDIDSGTEFGVTLGTPIGMLVLNDDQKP
MGSSFGKNFRVTTFGESHGGAVGVILDGCPPKLKINIDLIDNELDRRPGDSKITTPRNEKDKVDIDSGTEFGVTLGTPIGMLVLNDDQKP
MSSYGAALRIHTFGESHGSAVGCIIDGLPPRLPLSVEDVDPDLNRRPGDGPLSTDRREKDRVNILSGVEDGYTLGTPLAMLVWNEDRRP
MSTYGTLLKVTSYGESHGKAIGCVIDGFLSNIEINFDLIDKDLDRRPDDSKLTSNRNEKDKLVILSGFDENKTLGTPITFLIYNEDIKK
MSTYGTLLKVTFGESHGKAIGCVIDGFLSNIEINFDLIDKQLDRRPNDSKLTSNRNEKDKLVILSGFDENKTLGTPITFLIYNEDIKK
MSTYGTLLKVTFGESHGKGIGCVIDGFLPNITINFEEIDKQLNRRKPNDSKLTSNRNEKDKLVILSGFDDNKTLGTPITFLIENEDVKK
MSTYGTLLKVTFGESHGKGIGCVIDGFLPNITINFEEIDKQLNRRKPNDSKLTSNRNEPDKLIVLSGFDDNKTLGTPITFLIKNDVKK
MSTYGTLLKVTFGESHGKGIGCVIDGFLPNITINFEEIDKQLNRRKPNDSKLTSNRNEPDKLIVLSGFDDNKTLGTPITFLIKNDVKK
MSTYGTLLKVTFGESHGKGIGCVIDGFLPNITINFEEIDKQLNRRKPNDSKLTSNRNEPDKLIVLSGFDDNKTLGTPITFLIKNDVKK
MSTYGTLLKVTFGESHGKGIGCVIDGFLPNITINFEEIDKQLNRRKPNDSKLTSNRNEPDKLIVLSGFDDNKTLGTPITFLIKNDVKK
MSTYGTLLKVTFFGESHGKAIGCVIDGFLPNITINFEEIDKQLNRRKPNDSKLTSNRNEPDKLIVLSGFDDNKTLGTPITFLIKNDVKK
MSTYGTLLKVTFFGESHGKAIGCVIDGFLPNITINFEEIDKQLNRRKPNDSKLTSNRNEPDKLIVLSGFDDNKTLGTPITFLIKNDVKK

- A weak signal between 2 sequences may be stronger in the context of multiple sequences
- May allow construction of Phylogenetic trees
- May assist in protein structure prediction

#### Approach 1:

Select one sequence as the pivot. Perform pairwise alignment between each sequence and the pivot.

Locally Optimal

# Multiple Sequence Analysis

- Multiple Sequence Alignment (MSA)
  - Extension of Pairwise Alignments
  - Scoring
  - Branch and Bound Pruning
- Phylogenetic Trees – Structure and How to Compute
- Combining Phylogenetic Trees and MSA – Progressive Alignment – Weighting
- CLUSTAL





- What about the simple extension from 2D?
- There are seven possible "endings":



 $2^k - 1$  endings for k sequences. Why?

$$i,j,k = \max \begin{cases} s_{i-1,j-1,k-1} + \delta(u_i, v_j, w_k) \\ s_{i-1,j-1,k} + \delta(u_i, v_j, -) \\ s_{i-1,j,k-1} + \delta(u_i, -, w_k) \\ s_{i,j-1,k-1} + \delta(-, v_j, w_k) \\ s_{i-1,j,k} + \delta(u_i, -, -) \\ s_{i,j,k-1} + \delta(-, -, w_k) \\ s_{i,j-1,k} + \delta(-, v_j, -) \\ s_{i,j,k-1} + \delta(-, -, w_k) \end{cases}$$

S

 $\delta(x,y,z)$  is an entry in the 3D scoring matrix

Time and space grow exponentially with number of sequences.

# Scoring

Sum-of-Pairs Scoring (SP):

$$S(\mathcal{A}) = \sum_{i=1}^{m-1} \sum_{j=i+1}^{m} S(\bar{s}^i, \bar{s}^j)$$

Idea: A good multiple alignment should contain good pairwise alignments

# Scoring

Weighted Sum-of-Pairs Scoring (SP):

$$S(\mathcal{A}) = \sum_{i=1}^{m-1} \sum_{j=i+1}^{m} S(\bar{s}^i, \bar{s}^j) w^i w^j$$

The presence of many similar scores should not dominate the alignment, so we should incorporate some *a priori* knowledge of the sequences (i.e., organism/ family).

The dynamic programming matrix is large, but we only want the best alignment, and most matrix elements are not on that path.

Can we 'direct' the search to avoid evaluating cells that are provably not on the best path?



- $S_{v} \quad \mbox{Score of the best path from start} \\ \mbox{to } v$
- $F_{v} \quad \mbox{Bound on the best path from } v \ \mbox{to the end}$
- K Score of best known alignment

What if:  $S_v + F_v < K$ 

#### ARSTVK, ASVK, ARTR

Let v = (3, 2, 2)

 $S_v\,$  is score of best alignment of: ARS, AS, AR

 $F_v$  is upper bound on score of aligning: **TVK**, **VK**, **TR** 

If  $S_v + F_v < K$  then mark v as dead-ending (aka prune v)

#### begin

```
v = h_0; P(v) = 0; push(v, Q) ; push start cell on the queue
 while Q \neq \emptyset
   pop(v, Q); S(v) = P(v)
  if S(v) + F(v,h_N) \ge K then
    for all forward neighbors w of v do in the right order
      if w \notin Q then
        push(w, Q); P(w) = S(v) + D(v,w)
      else
       P(w) = \max(P(w), S(v) + D(v,w))
      end
    end
   end
 end
end
```

We know the alignment score is:

$$S(\mathcal{A}) \le \sum_{k=1}^{m-1} \sum_{l=k+1}^{m} S(s^k, s^l)$$

$$\frac{\text{Observation:}}{S(\mathcal{A}) = \sum_{i=1}^{m-1} \sum_{j=i+1}^{m} S(\bar{s}^i, \bar{s}^j)}$$
$$S(\bar{s}^i, \bar{s}^j) \le S(s^i, s^j)$$

So our bound can be:

$$F_{v} = \sum_{k=1}^{m-1} \sum_{l=k+1}^{m} S(s_{v_{k}+1...n_{k}}^{k}, s_{v_{l}+1...n_{l}}^{l})$$

Runtime for computing F (using dynamic programming):

$$O(n^2m^2)$$



# Phylogenetic Trees

- Goals
  - Reconstruct genealogical ties (the topology of the tree)
  - Estimate time of divergence (last common ancestor)

#### Given *m* original sequences

- -Tree should have *m* leaf nodes
- -Tree is rooted or unrooted
- -Set of internal nodes
  - Rooted tree: 2 children, one parent
  - Unrooted tree: 3 connected nodes
- -Internal nodes may have explicit sequence
- -Edges may carry a **weight** (representing *distance*)
- -Additive tree
  - Distance between two nodes is sum of edge distances
- -Ultrametric tree
  - An Additive tree where distance to common ancestors are equal

# Phylogenetic Trees

#### Additive Trees

- -Distance between two nodes is sum of edge distances
- -It's possible to construct an additive tree iff for any 4 nodes they can be labeled:  $D \rightarrow D \rightarrow D \rightarrow D \rightarrow D$

$$D_{i,j} + D_{k,l} = D_{i,k} + D_{j,l} \ge D_{i,l} + D_{j,k}$$

![](_page_16_Figure_5.jpeg)

- Ultrametric Trees
  - An Additive tree where distance to common ancestors are equal
  - It's possible to construct an ultrametric tree iff for every triple i, j, k

$$D_{i,j} \le \max(D_{i,k}, D_{k,j})$$

a = ba + c = b + c = d

d

![](_page_16_Picture_10.jpeg)

3 points arranged as an isosceles triangle with the odd side shorter than the other two

# **Constructing Phylo Trees**

#### Three Main Methods

#### -Maximum Parsimony

- Input: Multiple alignment
- Method: Find tree that minimizes number of evolutionary changes

#### - Maximum Likelihood

- Input: Multiple alignment
- Method: Identify statistically most likely tree (use substitution probs)

#### Distance-Based Methods

- Input: Multiple alignment
- Method: Use pairwise alignment distances, connect similar sequences

### **Distance-Based Methods**

Neighbor-Joining Method (NJ)

![](_page_18_Figure_2.jpeg)

- Start with all nodes arranged in a star topology
  - This topology has minimum number of edges
- Tree is successively changed by adding edges until an unrooted tree is constructed
- The two elements (a neighbor pair) that would create the overall smallest total edge length, are 'merged'

### Heuristic Alignment Progressive Alignment

- Let's return to Multiple Alignments
- Consider constructing an alignment 'pairwise'

```
C = \emptyset
for i = 1 to m do
C = C \bigcup \{\{s^i\}\}\}
end
for i = 1 to m - 1 do
choose two alignments A_p, A_q from C;
C = C - \{A_p, A_q\}A_r = align(A_p, A_q)C = C \bigcup \{A_r\}
end
```

Merge order is important

Method of aligning two sequences, an alignment and a sequence, or two alignments

# **Guiding Progressive Alignment**

**Phylogenetic Trees** 

![](_page_20_Picture_2.jpeg)

Multiple Seq Alignment

- Observation: Deterministic Multiple Seq Alignment is difficult
- Observation: Result of heuristic alignment very much depends on alignment ordering
- Observation: Multiple Sequence Alignment is supposed to reflect evolutionary relationships between sequences
- Idea: Use phylogenetic tree to guide alignment

![](_page_20_Figure_8.jpeg)

# **Guiding Progressive Alignment**

**Phylogenetic Trees** 

![](_page_21_Picture_2.jpeg)

Multiple Seq Alignment

- **Observation**: Deterministic Multiple Seq Alignment is difficult
- Observation: Result of heuristic alignment very much depends on alignment ordering
- Observation: Multiple Sequence Alignment is supposed to reflect evolutionary relationships between sequences
- Idea: Use phylogenetic tree to guide alignment

#### **Hierarchical Clustering**

Select 2 subset alignments using the Average Linkage Model Score is average score between all pairs of sequences (one from each alignment) Similar to PGMA

Align 2 subset alignments using Complete Alignment

#### Linear Clustering

Special case where one member being aligned must be a sequence

![](_page_22_Picture_0.jpeg)

![](_page_22_Picture_1.jpeg)

- Widely used Multiple Sequence Alignment Program
  - 1. Calculate the static pairwise similarity scores for the sequences
  - 2. Construct a guide tree by the use of the pairwise scores (NJ method)
  - 3. Calculate sequence weights, using the guide tree
  - 4. Perform a progressive alignment, guided by the tree

# Profiles

- Consider you have multiple genes from the same family and you want to find more.
  - -Pairwise alignment?
    - Slow, may miss weak signals
  - Multiple sequence alignment?
    - Slow, may be better with weak signals
- Other Alternatives?
- Profiles: A statistically grounded model describing the properties of the member sequences
  - Motifs
  - Position-Specific Scoring Matrices (PSSMs)
  - Profiles
  - Hidden Markov Model Profiles

# Constructing a Profile

- Let **Prof**<sub>*ra*</sub> be the score of aligning amino acid *a* (of the sequence) to the position *r* (of the profile)
- When computing Prof<sub>ra</sub> consider
  - Observed amino acids at position r
  - -Number of observations at position r
  - The similarity of a to the residues seen at position r
  - The background distribution of amino acids
  - -The weight of each sequence (as in MSA)
  - –Number of gaps at position r

$$\operatorname{Prof}_{ra} = \sum_{b \in \mathcal{M}} R_{ba} V_{rb} \qquad \text{where} \qquad$$

 $\mathcal{M}$ Alphabet of Amino Acids $R_{ba}$ Scoring Matrix $V_{rb}$ Weight of b at position r

# Constructing a Profile

• Consider  $V_{rb}$ - Ideally  $V_{rb} = 0$  when  $T_{rb} = 0$  $V_{rb} = 1$  when  $T_{rb} = m_r$ num residues (not gaps) at position r of alignment

-Can interpolate in various ways, or weight:

$$V_{rb} = \frac{\sum_{i=1}^{m_r} w_i \delta_b}{\sum_{i=1}^{m_r} w_i} \qquad \qquad \delta_b = 1 \text{ if } \bar{s}_r^i = b \\ \delta_b = 0 \text{ if } \bar{s}_r^i \neq b$$

#### • Gaps

 If *sufficient* number of gaps in position of MSA, remove that position from profile

### **PSI-BLAST**

- Position Specific Iterated BLAST (PSI-BLAST
- Idea
  - **Perform BLAST** with original query q
  - -Construct profile P from sequences aligned to q (pivot)
  - **Perform another BLAST** with P as new query
  - -Construct new profile P' from sequences
  - Repeat until convergence or max iterations
- Capable of identifying more distant homologs
- Useful in structure prediction

# Profile HMM

- Consider an HMM, constructed to emit sequences similar to some known profile
- Such an HMM might have states corresponding to insertions, matches, and deletions

![](_page_27_Figure_3.jpeg)

### Profile HMM

![](_page_28_Figure_1.jpeg)

Consider query sequence and best profile alignment

## Profile HMM

- Best path through HMM is found via DP (what were you expecting?)
  - Viterbi algorithm, recurrence is:

$$D_{i+1,j} = \max_{0 \le k \le j-1} (D_{i,k} P(k,j) P(q_{i+1}|j))$$

- •Note that for Profile HMMs, most transition probabilities are zero.
- •This leads to efficient implementations.
- •One may be interested in the sum of probabilities over all paths capable of generating the query sequence.

### **Profile Databases**

Many protein family databases with signatures for each family PFam, PROSITE, PRINTS, BLOCKS, ...

- •When searching a query against the database, compute the sequence emission probability for each family's HMM
- Report those above a threshold

![](_page_30_Picture_4.jpeg)