1/23/20

2. Homework (grad) Due 1/30/20 before class

1. Line segment intersection (5 points)

Explain how you can use one or more orientation tests to test if two line segments \overline{ab} and \overline{cd} intersect, where $a, b, c, d \in \mathbb{R}^2$. (Hint: Case analysis. Draw pictures of examples, and determine important configurations of a, b, c, d.)

2. Lower bounds (9 points)

Consider the following problems:

SORTING: Given a set $X = \{x_1, \ldots, x_n\}$ of *n* numbers, output the same numbers in non-decreasing order.

ELEMENT UNIQUENESS: Given a set $X = \{x_1, \ldots, x_n\}$ of *n* numbers, are there i, j, with $i \neq j$, such that $x_i = x_j$?

CLOSEST PAIR: Given a point set $P = \{p_1, \ldots, p_n\} \in \mathbb{R}^2$, output the closest pair of points in P.

ALL NEAREST NEIGHBORS: Given a point set $P = \{p_1, \ldots, p_n\} \in \mathbb{R}^2$. Compute for each point in P its *nearest neighbor* in P (i.e., point at minimum distance).

- (a) Prove a lower bound of $\Omega(n \log n)$ for SORTING, by reducing from ELEMENT UNIQUENESS (i.e., by using the knowledge that ELEMENT UNIQUENESS has a lower bound of $\Omega(n \log n)$).
- (b) Prove a lower bound of $\Omega(n \log n)$ for CLOSEST PAIR by reducing from an appropriate problem.
- (c) Prove a lower bound of $\Omega(n \log n)$ for ALL NEAREST NEIGHBORS by reducing from an appropriate problem.

3. Visible Segments Sweep (10 points)

Let S be a set of n disjoint line segments in the plane, and let p be a point not on any of the line segments of S. We say that the point p sees a line segment s if there is a point $q \in s$ such that the segment pq does not intersect any other line segment of S. We wish to determine all line segments of S that p can see.

Give an $O(n \log n)$ time algorithm for this problem that uses a rotating half-line with its endpoint at p to sweep the plane. You do not have to give pseudocode but you should explain all the necessary components of the sweep.

