## CMPS 2200 Introduction to Algorithms - Fall 15

9/22/15

## 4. Homework

Due $\mathbf{1 0} / \mathbf{1} / \mathbf{1 5}$ at the beginning of the lab

## 1. Recursion tree (8 points)

For the recurrences below, use the recursion tree method to find a good guess of what they could solve to asymptotically. Assume $T(1)=1$.
(a) (4 points) $T(n)=4 T(n / 3)+n^{2}$
(b) (4 points) $T(n)=8 T(n / 2)+n^{3}$
2. 3-way mergesort (4 points)

Consider the following variant of mergesort, where the first call is 3 wayMergesort ( $0, \mathrm{n}-1, \mathrm{~A}$ ) to sort the array $A[0 . . n-1]$.

```
int 3wayMergesort(int i, int j, int[] A){
    // Sort A[i..j]
    if(j-i<=1)
        return;
    int l = (j-i)/3;
    3wayMergesort(i,i+l, A);
    3wayMergesort(i+l+1,i+2*l,A);
    3wayMergesort(i+2*l+1,j,A);
    merge(i,i+l+1,i+2*l+1); // Merges all three sub-arrays in linear time
}
```

Set up a runtime recurrence for 3-way mergesort above. Then solve the recurrence using the method of your choice. What is the runtime of 3wayMergesort?

## 3. Divide-and-conquer (5 points)

Let $A$ be a sorted array of $n$ distinct integers. Give pseudo-code for a divide-andconquer algorithm that decides whether there is an index $i$ such that $A[i]=i$. Your algorithm should run in $O(\log n)$ time. Give a runtime recurrence for your algorithm and argue why the runtime is $O(\log n)$. (Hint: How many subproblems are you allowed to recurse into, in order to get achieve the required runtime?Also, remember that the array contains distinct integers.)

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4. Strassen's Algorithm (4 points)

Apply Strassen's algorithm to compute

$$
\left(\begin{array}{llll}
0 & 2 & 1 & 2 \\
3 & 1 & 0 & 1 \\
1 & 0 & 5 & 2 \\
2 & 1 & 1 & 1
\end{array}\right) \cdot\left(\begin{array}{llll}
1 & 1 & 1 & 1 \\
2 & 2 & 0 & 0 \\
1 & 3 & 1 & 2 \\
3 & 4 & 1 & 2
\end{array}\right)
$$

The recursion should exit with the base case $n=1$, i.e., $2 \times 2$ matrices should recursively be computed using Strassen's algorithm. In order to save you some work, you may assume that the following is a partial solution and you only have to fill in the missing values by using Strassen's algorithm:

$$
\left(\begin{array}{cccc}
11 & 15 & 3 & 6 \\
8 & 9 & 4 & 5 \\
& & 8 & 15 \\
& & 4 & 6
\end{array}\right)
$$

