

CMPS/MATH 2170 – Fall 2013

Graphs

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Graphs (review)

Definition. A *directed graph (digraph)* $G = (V, E)$ is an ordered pair consisting of

- a set V of *vertices* (singular: *vertex*),
- a set $E \subseteq V \times V$ of *edges*.

In an *undirected graph* $G = (V, E)$, the edge set E consists of *unordered* pairs of vertices.

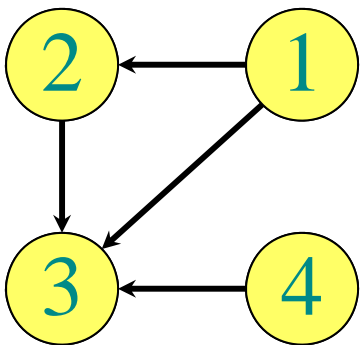
In either case, we have $|E| = O(|V|^2)$.

Moreover, if G is connected, then $|E| \geq |V| - 1$.

Adjacency-matrix representation

The *adjacency matrix* of a graph $G = (V, E)$, where $V = \{1, 2, \dots, n\}$, is the matrix $A[1 \dots n, 1 \dots n]$ given by

$$A[i, j] = \begin{cases} 1 & \text{if } (i, j) \in E, \\ 0 & \text{if } (i, j) \notin E. \end{cases}$$

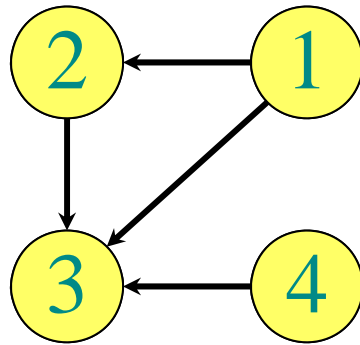


| A | 1 | 2 | 3 | 4 |
|---|---|---|---|---|
| 1 | 0 | 1 | 1 | 0 |
| 2 | 0 | 0 | 1 | 0 |
| 3 | 0 | 0 | 0 | 0 |
| 4 | 0 | 0 | 1 | 0 |

$\Theta(|V|^2)$ storage
 \Rightarrow *dense*
representation.

Adjacency-list representation

An *adjacency list* of a vertex $v \in V$ is the list $Adj[v]$ of vertices adjacent to v .



$$Adj[1] = \{2, 3\}$$

$$Adj[2] = \{3\}$$

$$Adj[3] = \{\}$$

$$Adj[4] = \{3\}$$

For undirected graphs, $|Adj[v]| = degree(v)$.

For digraphs, $|Adj[v]| = out-degree(v)$.

Adjacency-list representation

Handshaking Lemma:

Every edge is counted twice

- For undirected graphs:

$$\sum_{v \in V} \text{degree}(v) = 2|E|$$

- For digraphs:

$$\sum_{v \in V} \text{in-degree}(v) + \sum_{v \in V} \text{out-degree}(v) = 2|E|$$

⇒ adjacency lists use $\Theta(|V| + |E|)$ storage

⇒ a *sparse* representation

Paths, Cycles, Connectivity

Let $G=(V,E)$ be a directed (or undirected) graph

- A **path** from v_1 to v_k in G is a sequence of vertices v_1, v_2, \dots, v_k such that $(v_i, v_{i+1}) \in E$ (or $\{v_i, v_{i+1}\} \in E$ if G is undirected) for all $i \in \{1, \dots, k-1\}$.
- A path is **simple** if all vertices in the path are distinct.
- A path v_1, v_2, \dots, v_k forms a **cycle** if $v_1 = v_k$.
- A graph with no cycles is **acyclic**.
 - An undirected acyclic connected graph is called a **tree**. (Trees do not have to have a root vertex specified.)
 - A directed acyclic graph is a **DAG**. (A DAG can have undirected cycles if the direction of the edges is not considered.)
- An undirected graph is **connected** if every pair of vertices is connected by a path. A directed graph is **strongly connected** if for every pair $u, v \in V$ there is a path from u to v and there is a path from v to u .
- The **(strongly) connected components** of a graph are the equivalence classes of vertices under this reachability relation.