Algorithm Analysis Fall 2013 Carola Wenk



Nearly every modern electronic device can be thought of as a computer that transforms input to the desired output.



Most computers are general-purpose: we can accomplish a number of different tasks on a single piece of hardware by changing the program being executed.



A program is just a realization (= implementation) of an abstract procedure, or algorithm, on a particular hardware platform (= computer). One algorithm can be used for a variety of applications.



We think of an algorithm as a sequence of actions to take input and produce output.

We assume a model of computation, usually an abstract instruction set (arithmetic operations, if-then, loops) and assign unit cost (= time) to them.

Describe an Algorithm

- 1. Define the problem (input, output)
- 2. Describe the algorithm (in words or in pseudo-code)
- 3. Proof of correctness (convince the reader of correctness)

Algorithm

Output

4. Analysis (runtime, space)

Describe an Algorithm:
Computing the minimum of *n* numbers
1. Define the problem (input, output)
Input: A list of *n* numbers.

Output: The value of the minimum number. [Other option: The index of the minimum number.]

Describe an Algorithm: Computing the minimum of *n* numbers

- 2. Describe the algorithm (in words or in pseudo-code)
- Loop through all numbers.
- Keep track of the minimum number seen so far.
- If current number is less than the minimum seen so far, update the minimum.

Describe an Algorithm: Computing the minimum of *n* numbers 2. Describe the algorithm (in words or in pseudo-code) def my_min(list):

min_so_far = list[0]
for i in range(1,len(list)):
 if list[i] < min_so_far:
 min_so_far = list[i]
 return min_so_far</pre>

Describe an Algorithm: Computing the minimum of n numbers 3. Proof of correctness (convince the reader of correctness) list 0 1 2 ... n-1 def my min(list): i-1.i min so far = list[0] for i in range(1,len(list)): At beginning of loop if list[i] < min so far: body: min so far = list[i] min so far = minimum return min so far of list[0]...list[i-1]

 Loop body updates min_so_far

Describe an Algorithm: Computing the minimum of *n* numbers 4. Analysis (runtime, space) Let n = len(list)Instructions: def my min(list): min so far = list[0] for i in range(1,len(list)): if list[i] < min so far: $\leq 3(n-1)$ min_so_far = list[i] return min so far Runtime: $\leq 2 + 3n - 3 = (3n - 1)$ Runtime linear in n Runtime O(*n*)

Slow algorithm to compute the minimum of *n* numbers

• Let n = len(list)

def my_min_slow(list): min_so_far = list[0] for j in range(1,len(list)+1): for i in range(0,j): if list[i] < min_so_far: min_so_far = list[i] return min_so_far 1 Instructions: 1 $\sum_{j=1}^{n} (3j+1)$ $\leq 3j$ $\sum_{j=1}^{n} (n(n+1)) + n$ $\leq 3j$

• Runtime: $\leq 2 + \frac{3}{2}n^2 + \frac{5}{2}n$ Runtime quadratic in nRuntime $O(n^2)$ Runtimes: Functions in input size *n* Runtime for my_min: f(n)=3n-1Runtime for my_min_slow: $g(n) = 2 + \frac{3}{2}n^2 + \frac{5}{2}n$

Which one is better? And for what values of n?

Asymptotic Runtime Analysis

To evaluate the abstract runtime of an algorithm, we want to know its <u>asymptotic behavior</u> as a function of the input size.

 \rightarrow How does the algorithm perform if the input grows larger and larger?



The growth rate of the running time allows us to compare and contrast two potential algorithms <u>before</u> implementing them.

Input Size

(Worst-Case) Asymptotic Runtime Analysis Usually, the abstract performance of an algorithm

depends on the actual input for any particular size n.

Which inputs should we use to characterize runtime?



We define algorithm performance as conservatively as possible, on the worst-case inputs.

"No matter what, my algorithm →takes at most c.*n* steps for an input size of *n*."

Input Size

(Worst-Case) Asymptotic Runtime Analysis Usually, the abstract performance of an algorithm

depends on the actual input for any particular size n.

Which inputs should we use to characterize runtime?



Input Size

We define algorithm performance as conservatively as possible, on the worst-case inputs.

Definition: $f(n) \in O(g(n))$ iff $\rightarrow \exists c, n_0 > 0 \forall n \ge n_0$: $f(n) \le c g(n)$