

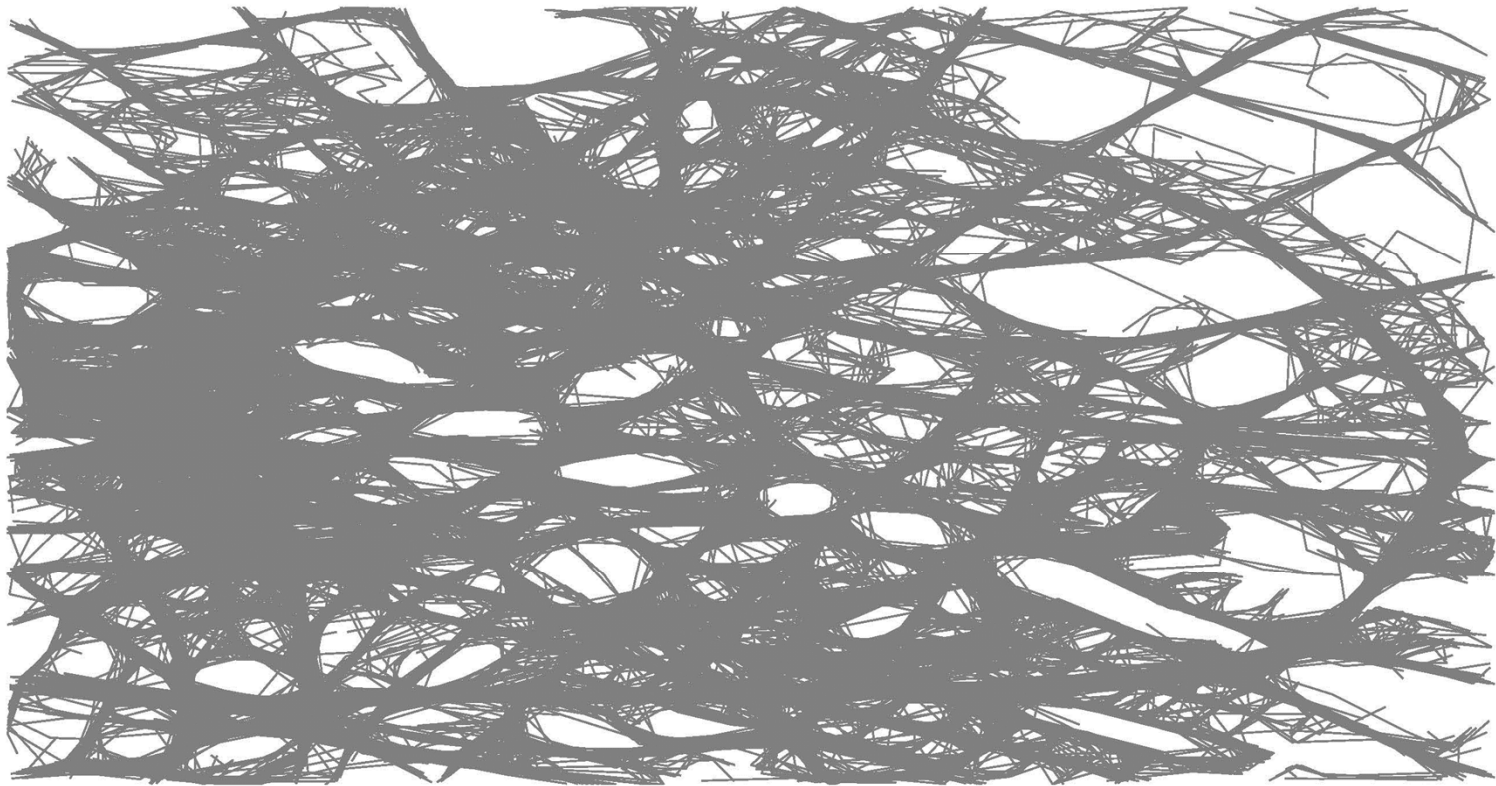


Constructing Road Maps from Trajectories

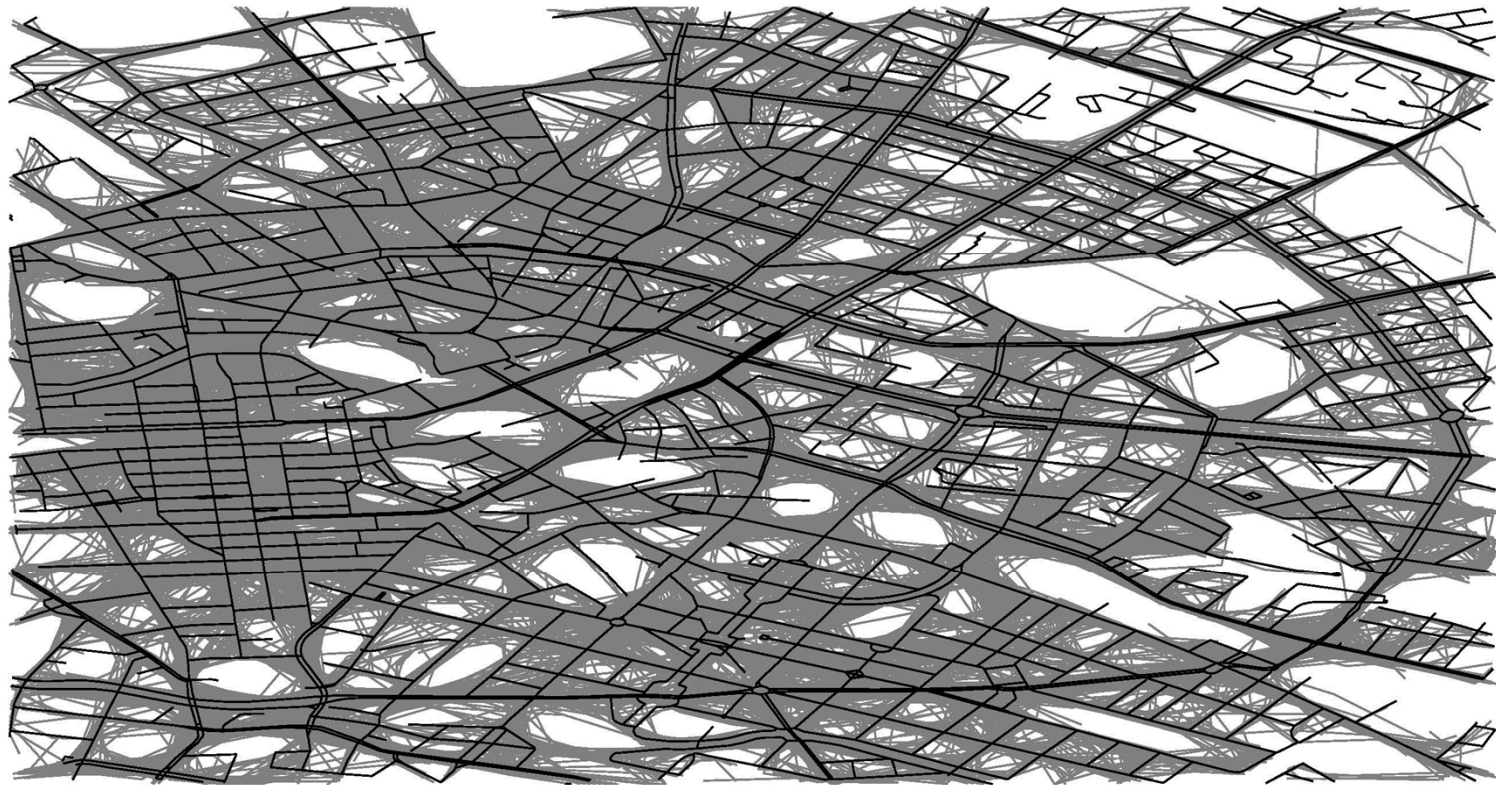
Carola Wenk

Department of Computer Science
Tulane University

GPS Trajectory Data



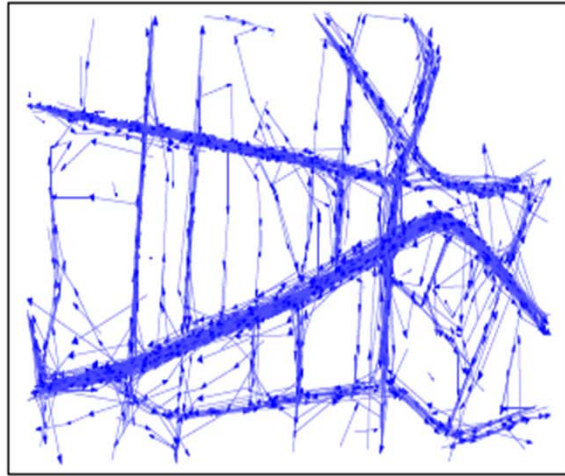
GPS Trajectory Data & Roadmap



⇒ Map Construction

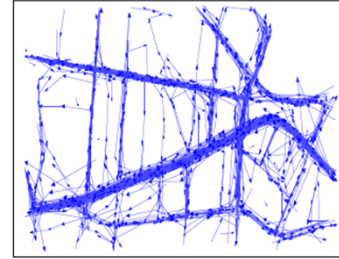
Map Construction

- Given a set of trajectories, compute the underlying road network



- Capturing constrained movement (explicit or implicit streets/routes, animal behavior)
- mapconstruction.org , openstreetmap.org
- Related problem: Map updates

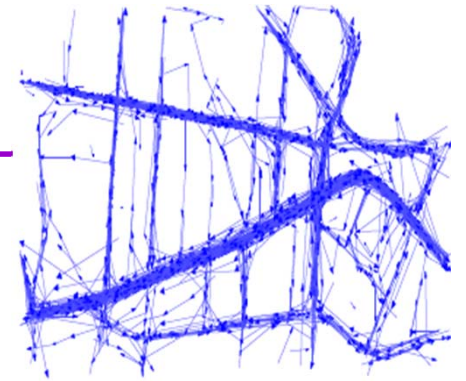
Map Construction



Geometric reconstruction problem:

- Given a set of movement-constrained trajectories, extract the underlying geometric graph structure
- Reconstruct a geometric domain that has been sampled with continuous curves that are subject to noise
 - ⇒ Sampling with organized data (trajectories) instead of point clouds
 - ⇒ Need to identify combinatorial information (edges, vertices), as well as geometric representation/embedding
 - ⇒ Clustering & how to represent an edge/street

Trajectories



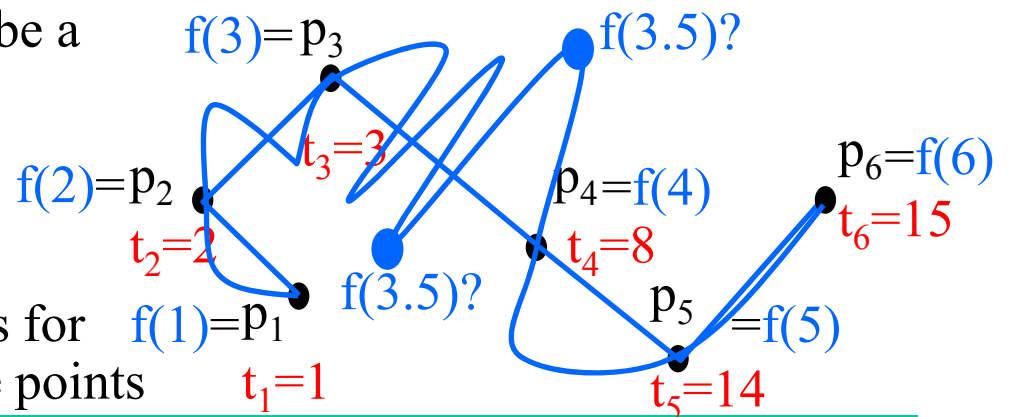
- A **trajectory** is a sequence of position samples: p_1, \dots, p_n
- Each p_i minimally consists of:
 - position measurement (e.g., (x,y)-coordinate)
 - time stamp
 - \Rightarrow e.g., $p_i = (x_i, y_i, t_i)$

- Such a trajectory is a finite sample of a **continuous curve** $f: [t_1, t_n] \rightarrow \mathbb{R}^2$

- For simplicity, f is often assumed to be a piecewise linear interpolation.

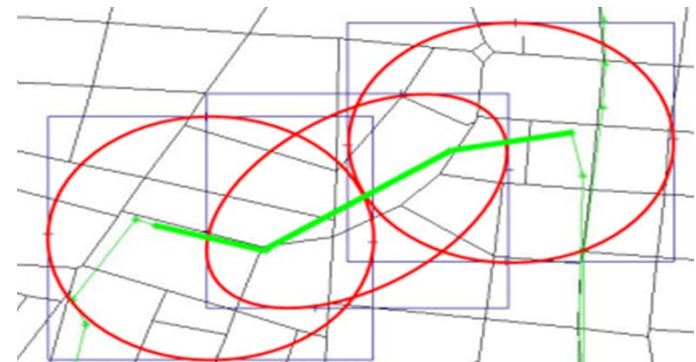
- But clearly there are many possible choices for f .

- There are also many possible choices for parameterizations in between sample points



Uncertainty and Error/Noise

- **Measurement error:** Usually modeled as Gaussian noise, or as an error-disk around each measurement point.
 - **Sampling error:**
 - Amounts to modeling the transition between two measurements
 - Simple transition model: Linear interpolation.
Common transition models in ecology: Brownian bridges, Levy walks
 - Simple region-based model: Buffers of fixed radius around each trajectory
- ⇒ Need **input model**:
E.g., chain of beads model for trajectories
- ⇒ What is a good **output model**?

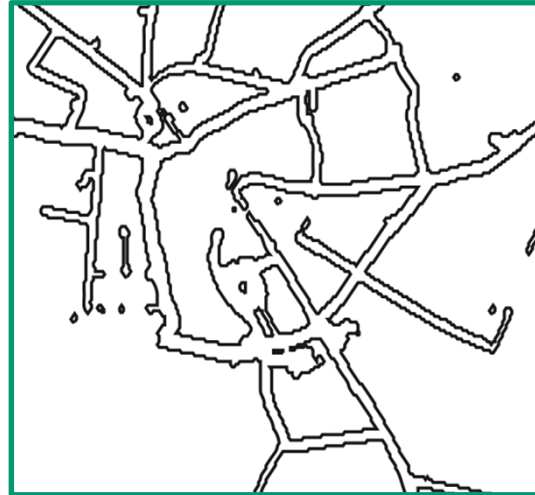


Map Construction: Some Results

- [DBH06]: Classical Kernel Density Estimation based method



Density of tracks



Contour



Center lines

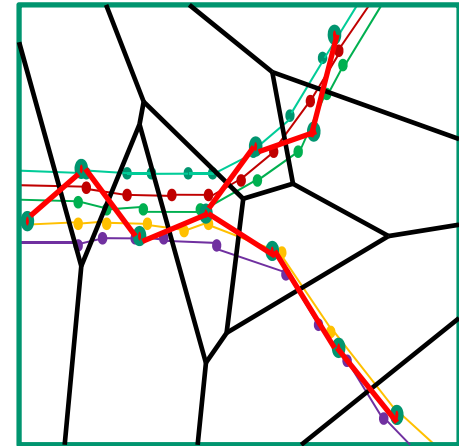
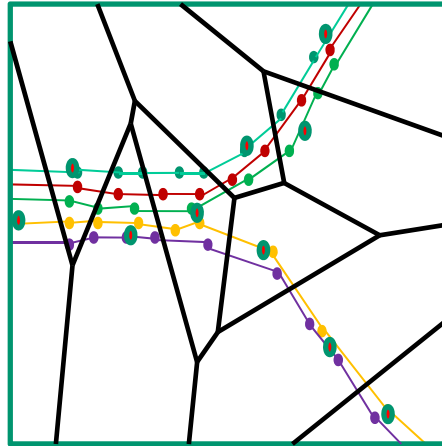
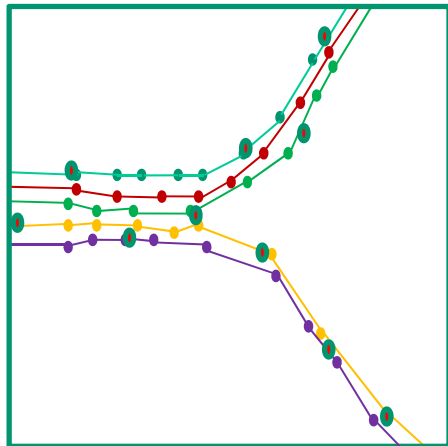
- [BE12]: Kernel Density Estimation based method; pipeline to first create scaffold then map-match trajectories.

[DBH06] J. Davies, A. Beresford, A. Hopper: Scalable, distributed, real-time map generation. IEEE Pervasive Comp. 5(4), 47-54, 2006.

[BE12] J. Biagioni, J. Eriksson, Map inference in the face of noise and disparity, 20th ACM SIGSPATIAL: 79-88, 2012

Map Construction: More Results

- [CGHS10]: First algorithm with quality guarantees. Subsamples trajectories, yielding a point cloud. Uses local neighborhood simplicial complexes. Reconstructs “good” portions of edges.



- [ACCGGM11]: Reconstruct metric graph from point cloud. Compute almost isometric space with lower complexity. Focus on combinatorial info and not embedding. Quality guarantees assume dense sampling.
- [GSBW11]: Topological approach on neighborhood complex. Uses Reeb graph to model skeleton graph (branching structure)

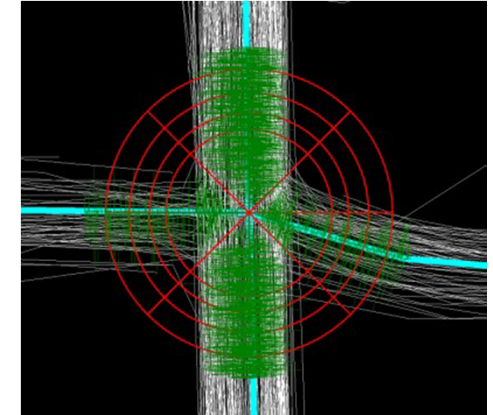
[ACCGGM11] M. Aanjaneya, F. Chazal, D. Chen, M. Glisse, L. Guibas, D. Morozov. Metric graph reconstruction..., SoCG, 2011.

[CGHS10] D. Chen, L. Guibas, J. Hershberger, J. Sun, Road network reconstruction for organizing paths, SODA, 2010.

[GSBW11] X. Ge, I. Safa, M. Belkin, Y. Wang, Data skeletonization via Reeb graphs, Conf. Neural Inf. Proc. Systems: 837-845, 2011.

Map Construction: Even More Results

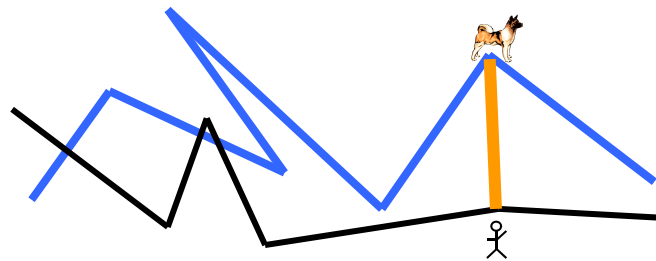
- [FK10]: First identify intersections (vertices) using a shape descriptor, then fill in edges.
- [KP12]: Detect intersections from turns and speed change, then fill in edges.



- [AW12]: Use trajectory information. Incrementally add one trajectory after another. Use partial Fréchet distance to identify new and existing portions. Use min-link algorithm to compute representative curve/edge.

[FK10] A. Fathi, J. Krumm, Detecting road intersections from GPS traces, Geographic Information Science, LNCS 6292: 56-69, 2010.
[KP12] S. Karagiorgou, D. Pfoser, On vehicle-tracking data-based road network generation, 20th ACM SIGSPATIAL: 89-98, 2002.
[BE12] J. Biagioni, J. Eriksson, Map inference in the face of noise and disparity, 20th ACM SIGSPATIAL: 79-88, 2012
[AW12] M. Ahmed, C. Wenk, Constructing Street Networks from GPS Trajectories, ESA: 60-71, 2012.

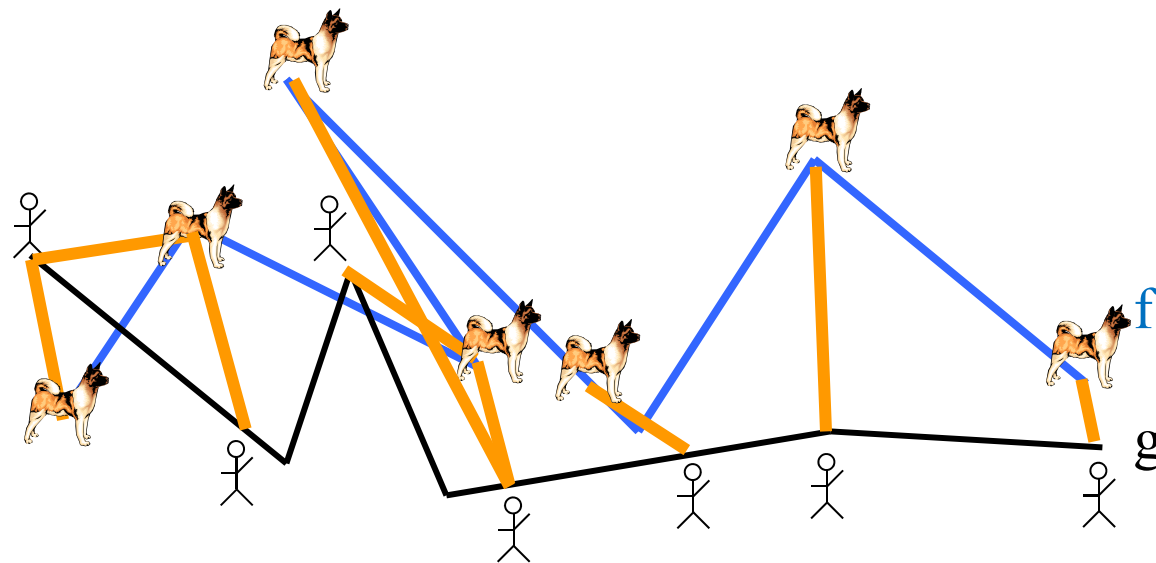
Fréchet Distance



Fréchet Distance for Curves

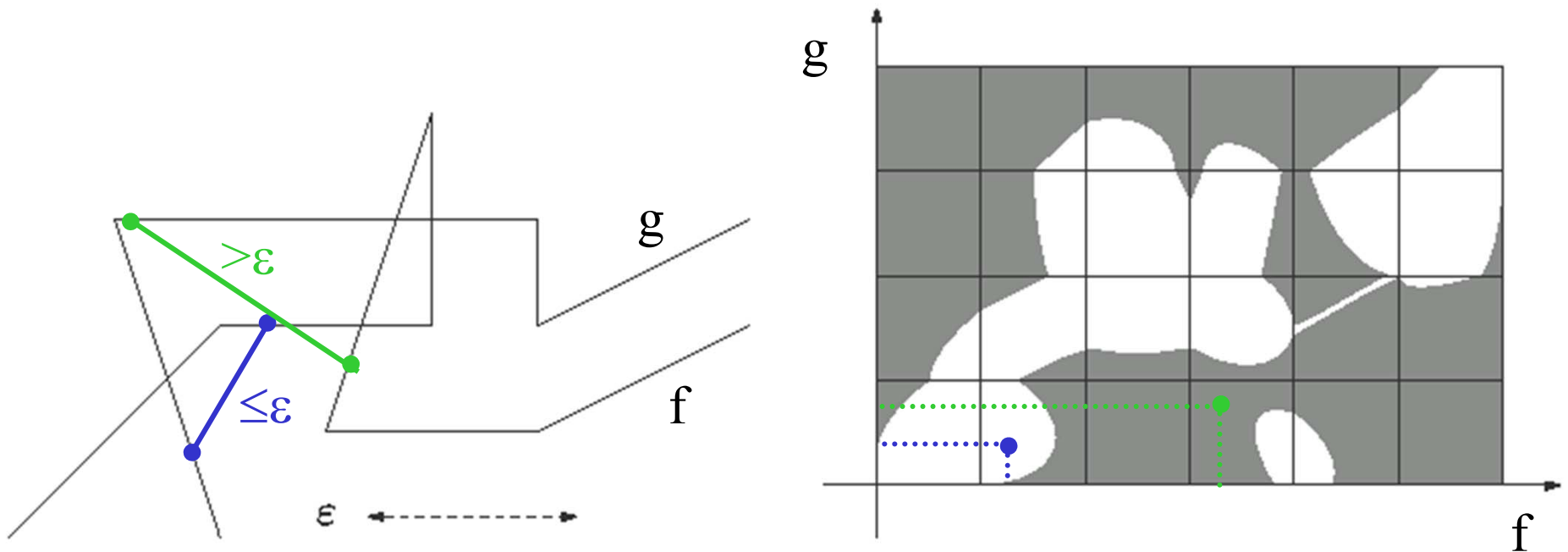
$$\delta_F(f,g) = \inf_{\alpha, \beta: [0,1] \rightarrow [0,1]} \max_{t \in [0,1]} \|f(\alpha(t)) - g(\beta(t))\|$$

where α and β range over continuous monotone increasing reparameterizations only.



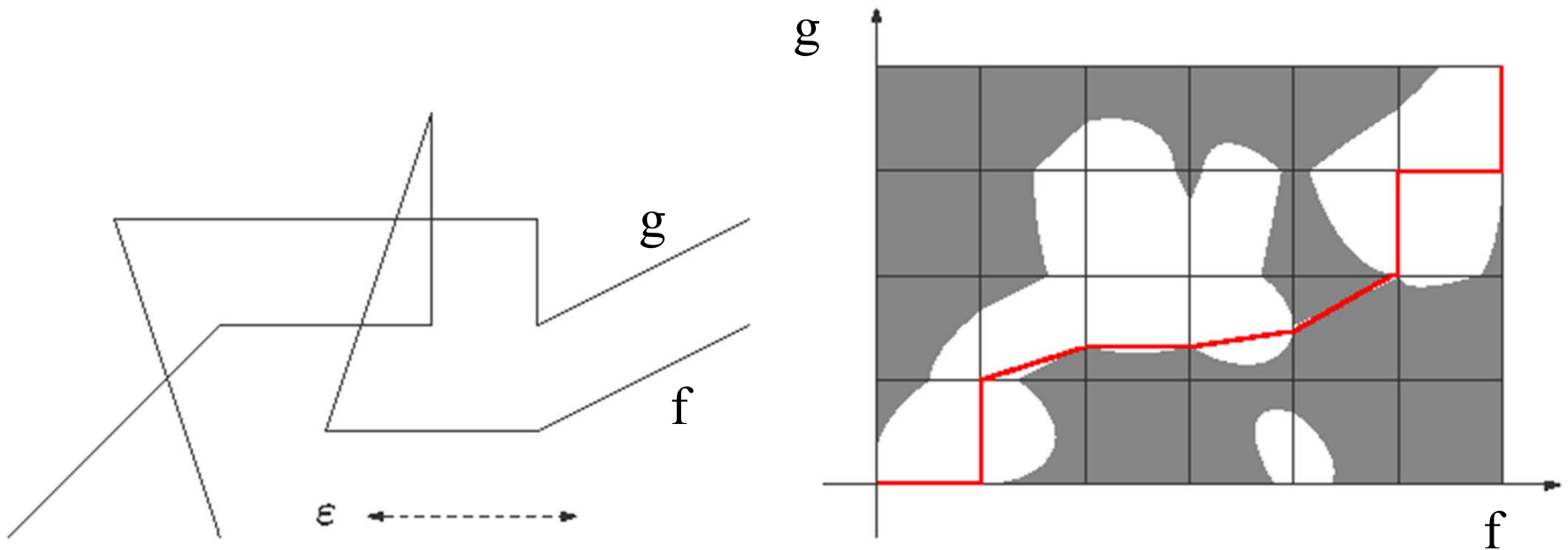
- Man and dog walk on one curve each
- They hold each other at a **leash**
- They are only allowed to go forward
- δ_F is the minimal possible leash length

Free Space Diagram



- Let $\epsilon > 0$ fixed (eventually solve decision problem)
- $F_\epsilon(f,g) = \{ (s,t) \in [0,1]^2 \mid \| f(s) - g(t) \| \leq \epsilon \}$ *white points*
free space of f and g

Free Space Diagram



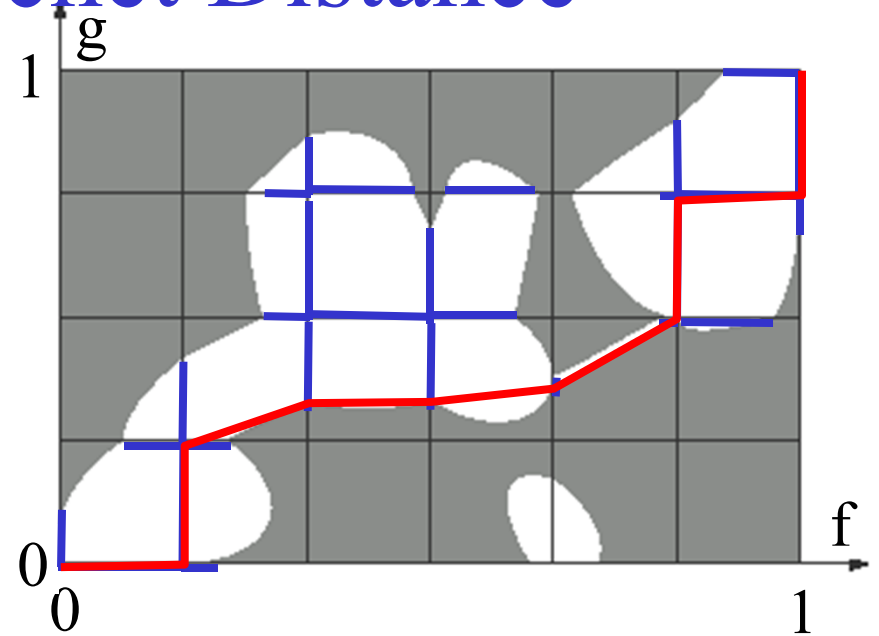
- Monotone path encodes reparametrizations of f and g
- $\delta_F(f,g) \leq \varepsilon$ iff there is a monotone path in the free space from $(0,0)$ to $(1,1)$
- Such a path can be computed using DP in $O(mn)$ time

Compute the Fréchet Distance

- **Solve the decision problem**

$\delta_F(f,g) \leq \varepsilon$ in $O(mn)$ time:

- Find monotone path using DP:
- On each cell boundary compute the interval of all points that are reachable by a monotone path from $(0,0)$
- Compute a **monotone path** by backtracking



- **Solve the optimization problem**

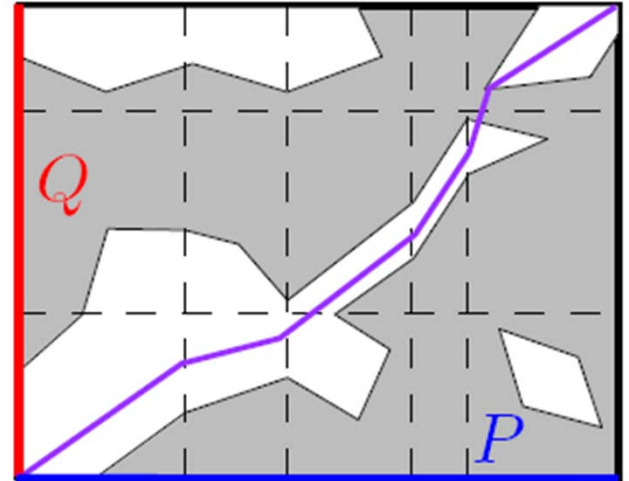
- In practice in $O(mn \log b)$ time with binary search and b-bit precision
- In $O(mn \log mn)$ time [AG95] using parametric search (using Cole's sorting trick)
- In $O(mn \log^2 mn)$ expected time [CW09] with randomized red/blue intersections

[AG95] H. Alt, M. Godau, Computing the Fréchet distance between two polygonal curves, *IJCGA* 5: 75-91, 1995.

[CW10] A.F. Cook IV, C. Wenk, Geodesic Fréchet Distance Inside a Simple Polygon, *ACM TALG* 7(1), 19 pages, 2010.

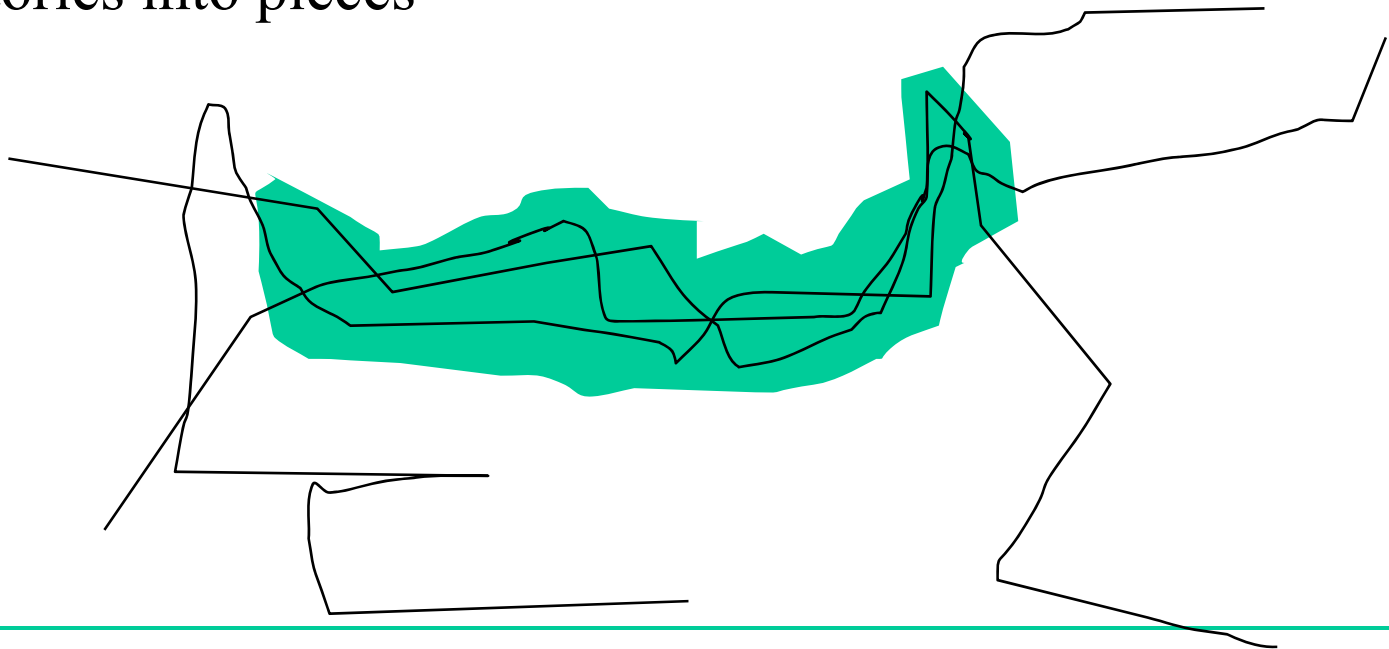
Partial Fréchet Distance

- For a given $\varepsilon > 0$, compute a monotone path in the free space diagram
 - that is allowed to pass through both white and black regions and
 - that maximizes the portion of the path within the white regions.
- Apply DP approach as before, but on each cell boundary maintain a function (instead of an interval). This function measures the maximum length of any monotone path from the lower left corner to the point on the boundary.
- For technical reasons the L_1 -distance is used to measure the Fréchet distance (hence the free space is polygonal)
- Runtime $O(n^3 \log n)$
- This partial distance identifies portions of the two curves that correspond to each other \Rightarrow helpful for subtrajectory clustering



Sub-Trajectory Clustering

- Find similar portions in trajectories
- Lots of algorithms for finding clusters in point sets
- Harder for trajectories since you need to figure out where to break the trajectories into pieces



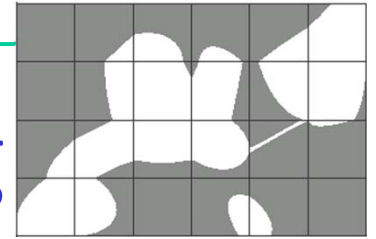
Sub-Trajectory Clustering

- A. Asahara, A. Sato, and K. Maruyama. Evaluation of trajectory clustering based on information criteria for human activity analysis. In 10th Int. Conf. on Mobile Data Management: Systems, Services and Middleware (MDM), pages 329-337, 2009.
 - K. Buchin, M. Buchin, J. Gudmundsson, M. Löffler, and J. Luo. Detecting commuting patterns by clustering subtrajectories. International Journal of Computational Geometry and Applications, special issue on 19th International Symposium on Algorithms and Computation (ISAAC), 2010.
 - K. Buchin, M. Buchin, M. van Kreveld, and J. Luo. Finding long and similar parts of trajectories. In 17th ACM SIGSPATIAL International Conference on Advances in Geographic Information Systems (ACM GIS), pages 296-305, 2010.
 - A. Dahlbom and L. Niklasson. Trajectory clustering for coastal surveillance. In 10th Int. Conf. on Information Fusion, pages 1-8, 2007.
 - J. Lee, J. Han, and K.-Y. Whang. Trajectory clustering: A partition-and-group framework. In Proc. ACM SIGMOD International Conference on Management of Data, pages 593-604, 2007.
-

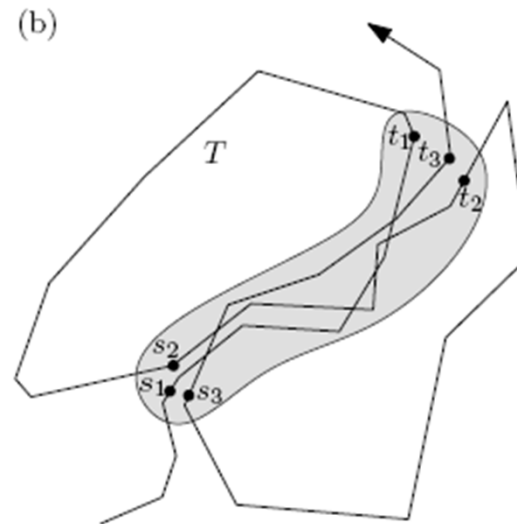
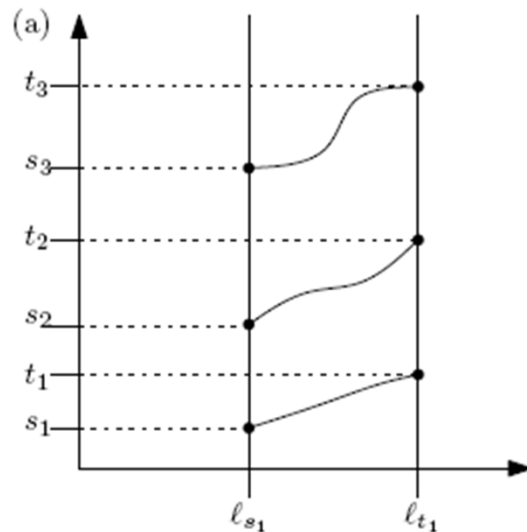
Sub-Trajectory Clustering

- X. Li, W. Hu, and W. Hu. A coarse-to-ne strategy for vehicle motion trajectory clustering. In 18th Int. Conf. on Pattern Recognition (ICPR), volume 1, pages 591-594, 2006.
 - Z. Li. Incremental clustering for trajectories. Master's thesis, University of Illinois at Urbana-Champaign, 2010.
 - Z. Li, J.-G. Lee, X. Li, and J. Han. Incremental clustering for trajectories. In Proc. 15th Int. Conf. Database Systems for Advanced Applications (DASFAA), pages 32-46, 2010.
 - T.W. Liao. Clustering of time series data - a survey. Pattern Recognition, 38:1857-1874, 2005.
 - Y. Zhang and D. Pi. A trajectory clustering algorithm based on symmetric neighborhood. In WRI World Congress on Computer Science and Information Engineering, volume 3, pages 640-645, 2009.
-

Fréchet-Based Clustering



- Given an input set of trajectories, append them all to form a single trajectory f
- Compute the free space diagram of f with itself (comparing f with f)



- Find clusters of monotone curve pieces in the (white) free space
⇒ Sweep free space from left to right, maintain data structure

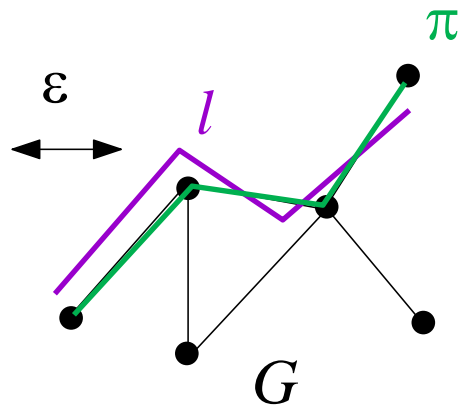
Map-Matching



Map Matching

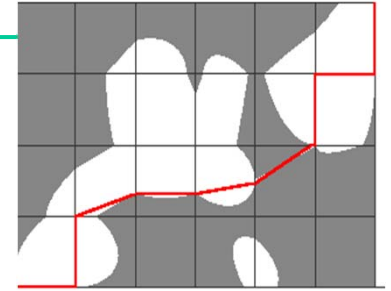
Given: A graph G , a curve l , and a distance parameter ε .

Task: Find a path π in G such that $\delta_F(l, \pi) \leq \varepsilon$



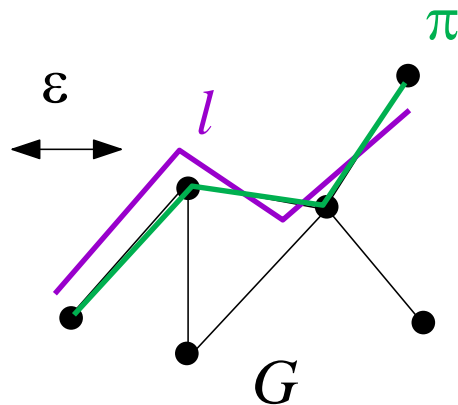
Application: GPS routing; use GPS data from vehicle fleets to build data base of current travel times

Subtask: Map Matching

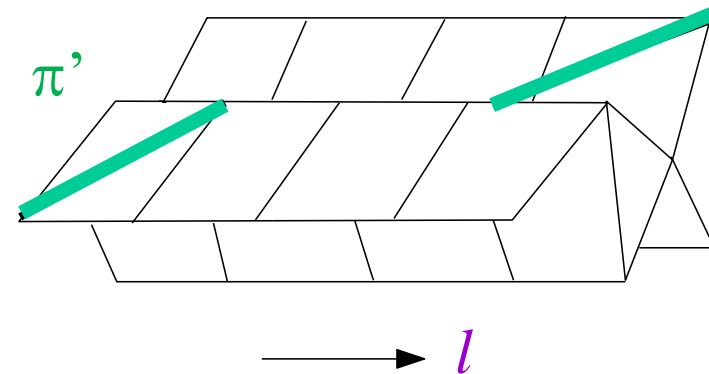


Given: A graph G , a curve l , and a distance parameter ε .

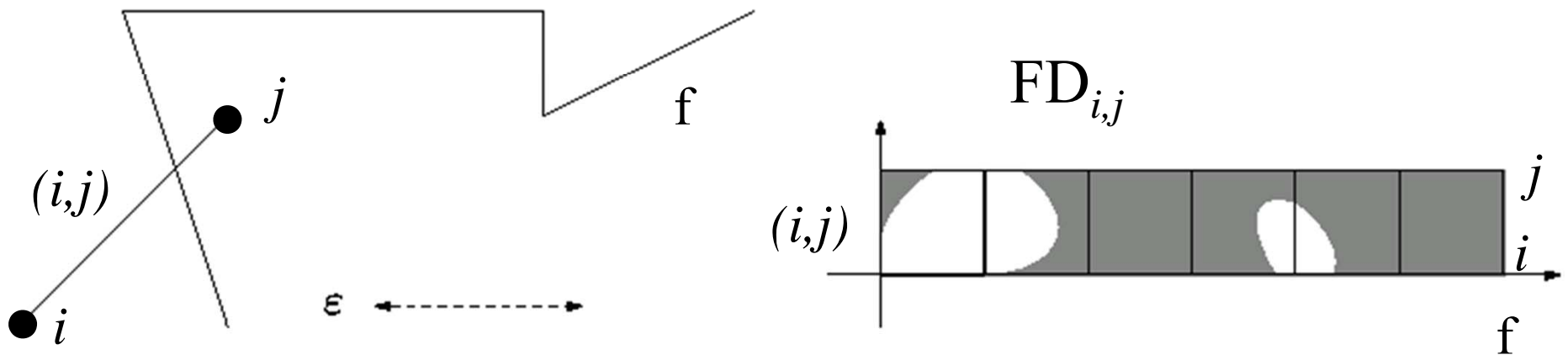
Task: Find a path π in G such that $\delta_F(l, \pi) \leq \varepsilon$



Compute free space surface.
And find path π' in it



Map Matching: Free Space Diagram



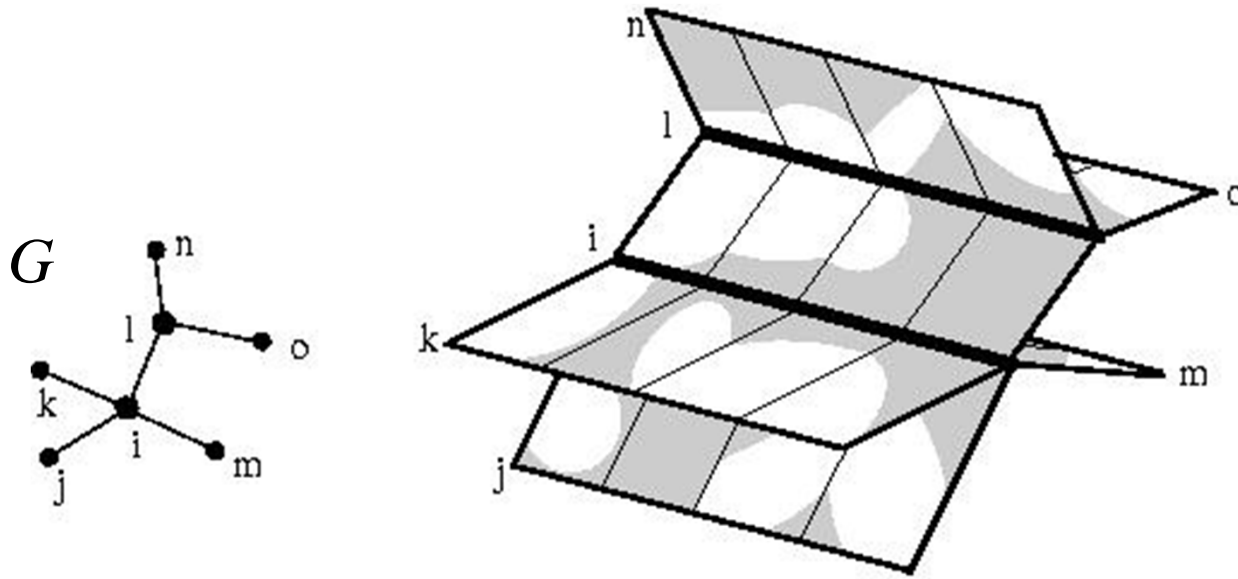
- For every edge (i,j) in G : $FD_{i,j} = FD(f, (i,j))$
- For every vertex i in G : sei $FD_i = FD(f, i)$ *1-dimensional*

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Free Space Surface



- Glue the free space diagrams $FD_{i,j}$ together according to adjacency information in G

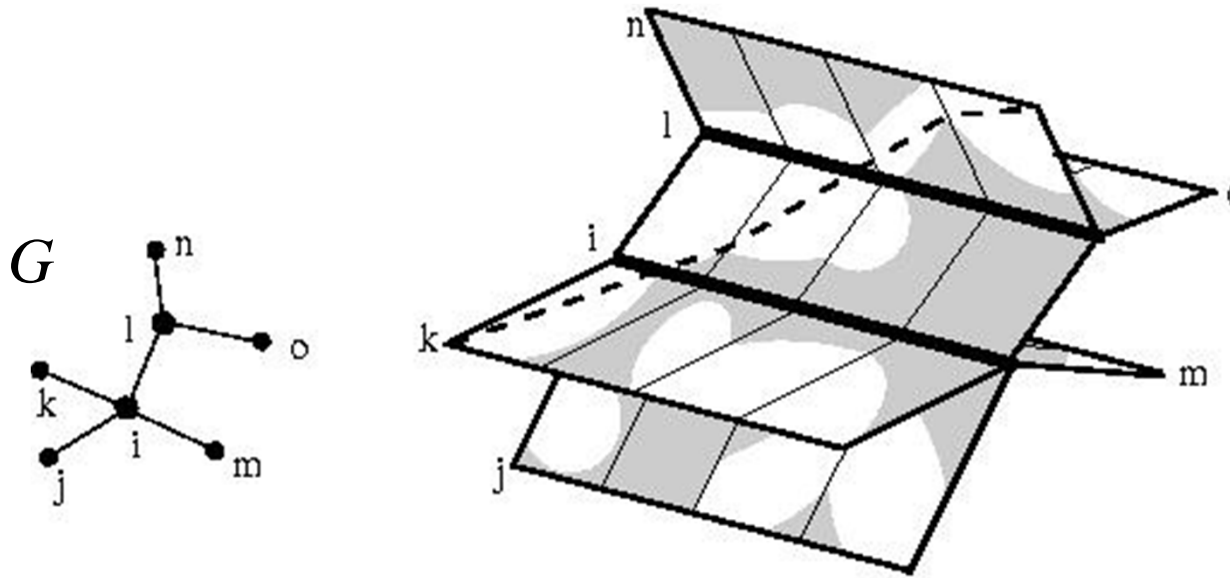
➔ **Free space surface of f and G**

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Free Space Surface



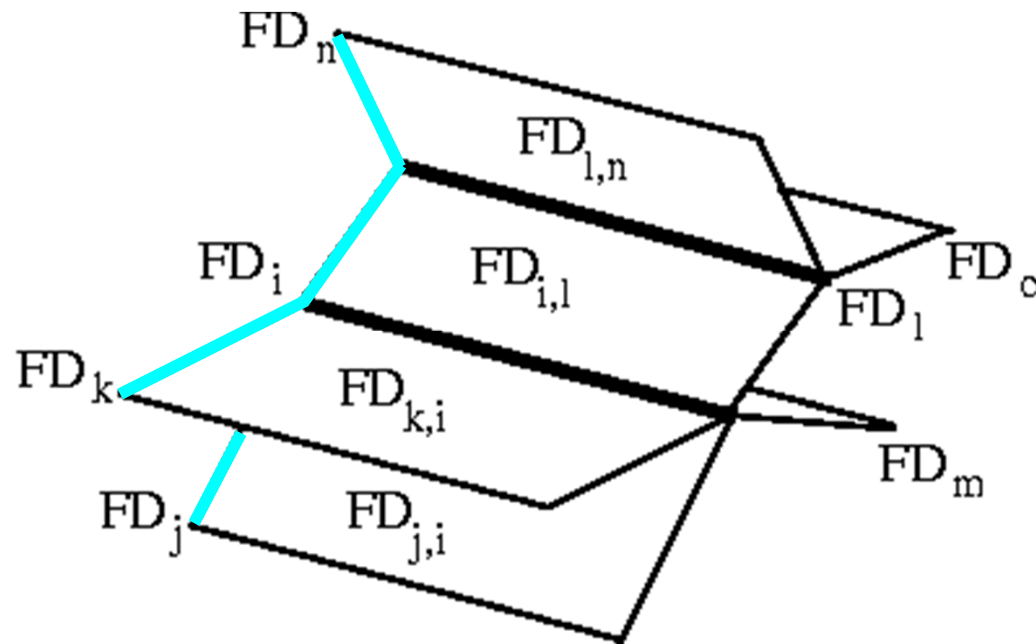
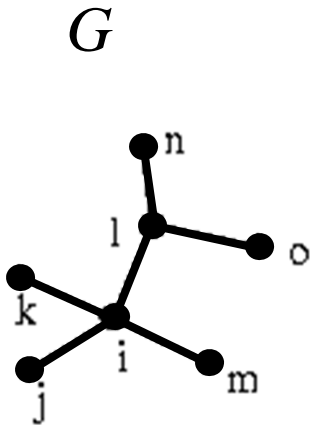
- **Task:** Find a **monotone path** in the free space surface that
 - starts in a lower left corner
 - and ends in an upper right corner

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Sweep



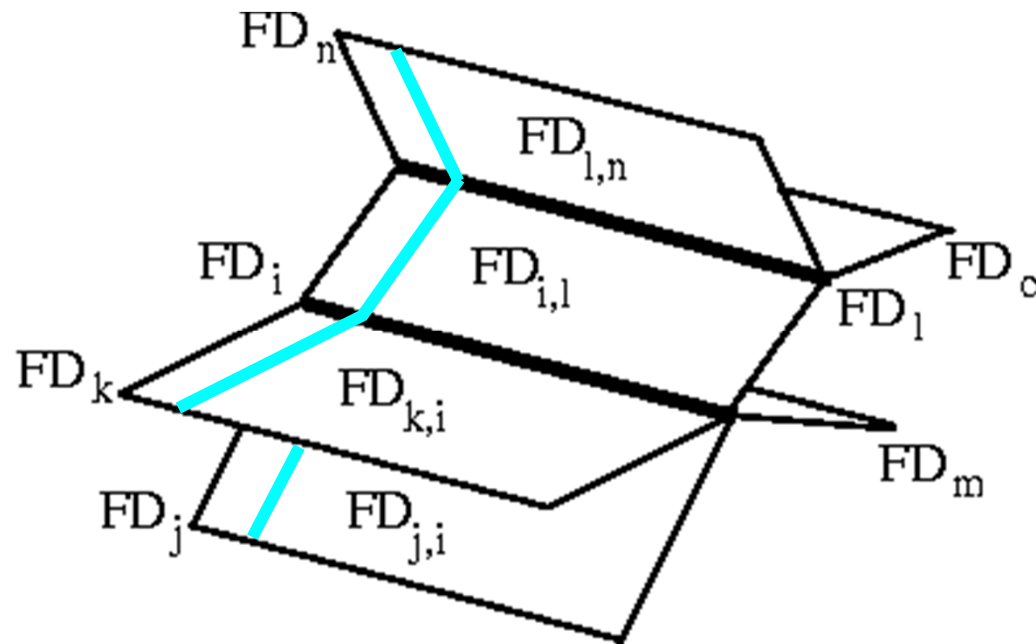
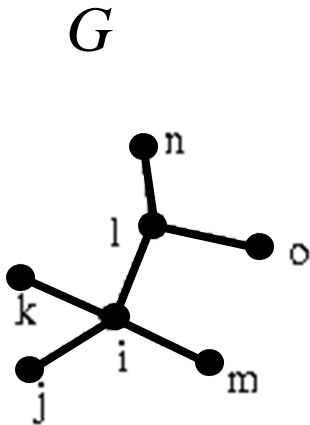
- Sweep all $FD_{i,j}$ with a **sweep line** from left to right

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Sweep



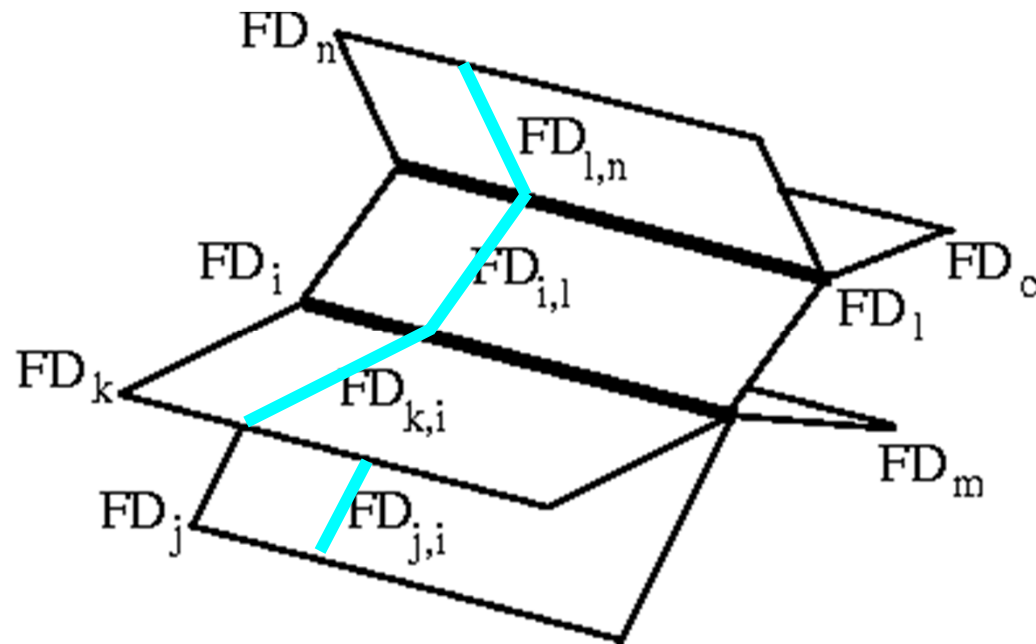
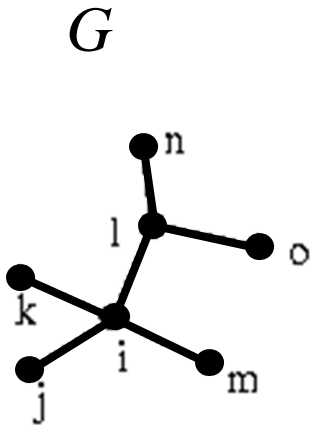
- Sweep all $FD_{i,j}$ with a **sweep line** from left to right

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Sweep



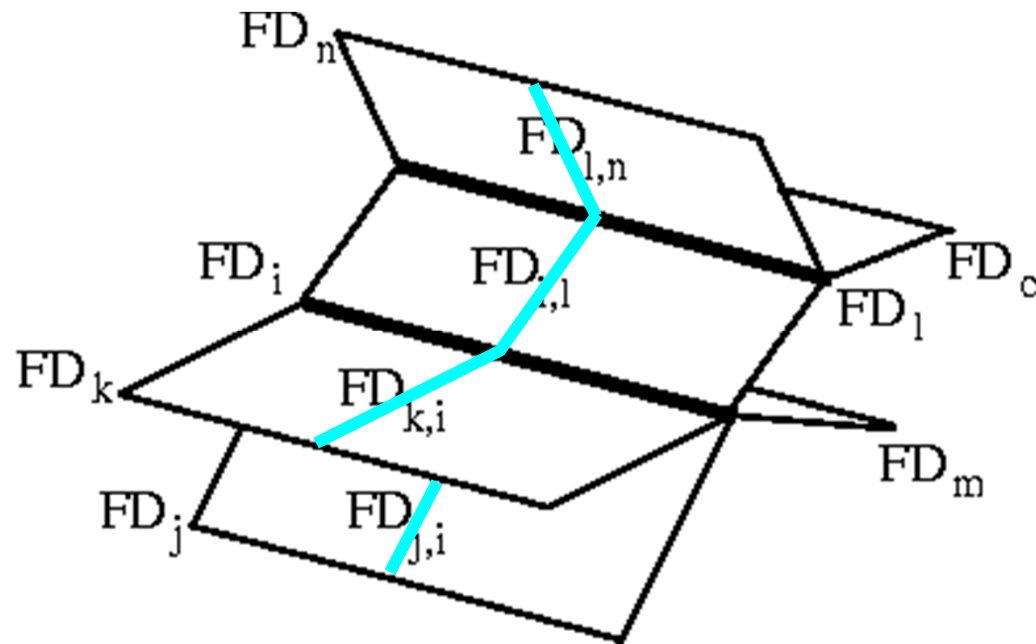
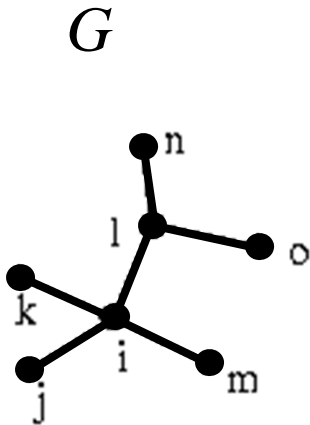
- Sweep all $FD_{i,j}$ with a **sweep line** from left to right

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Sweep



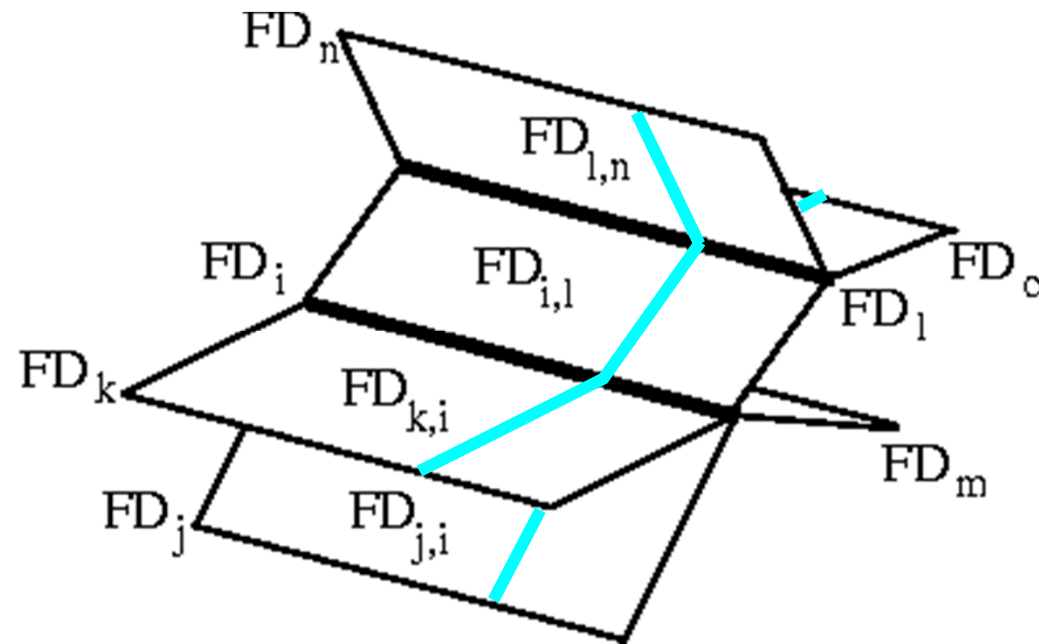
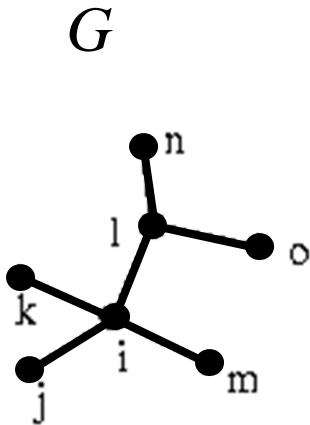
- Sweep all $FD_{i,j}$ with a **sweep line** from left to right

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Sweep



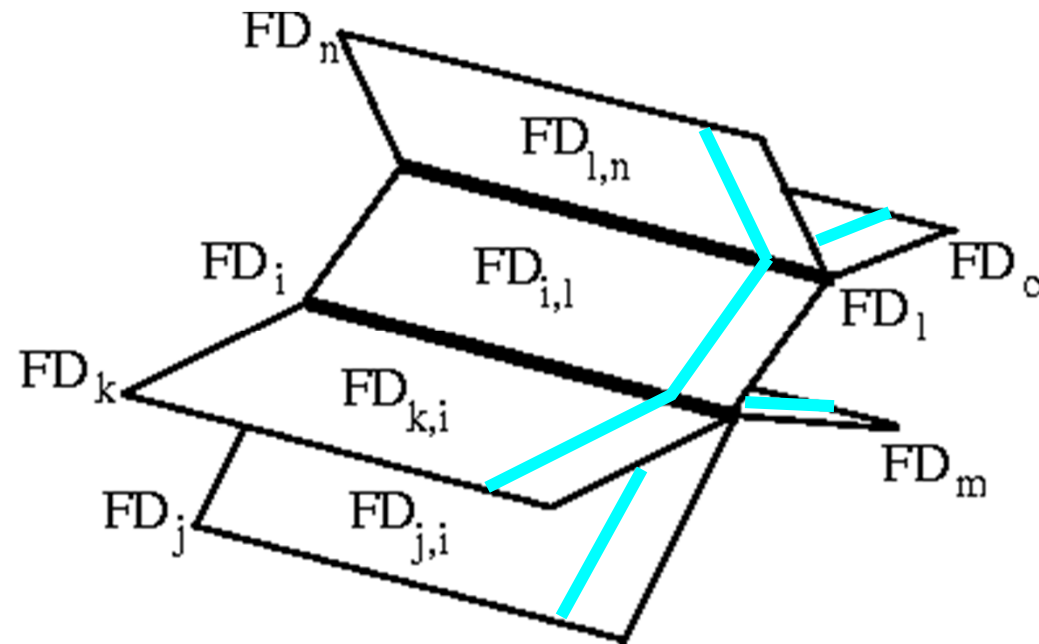
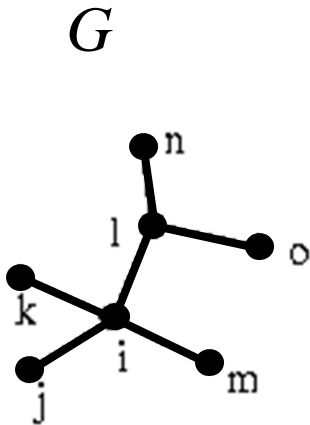
- Sweep all $FD_{i,j}$ with a **sweep line** from left to right

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Sweep



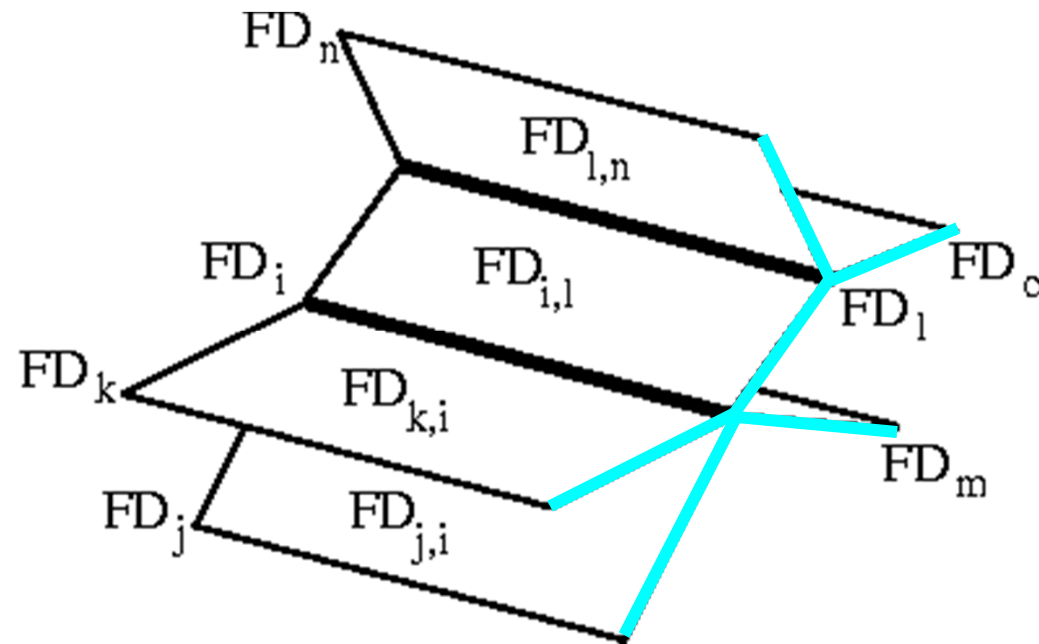
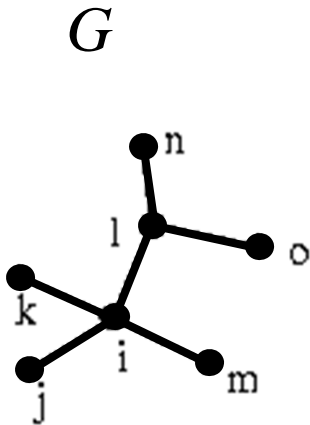
- Sweep all $FD_{i,j}$ with a **sweep line** from left to right

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Sweep



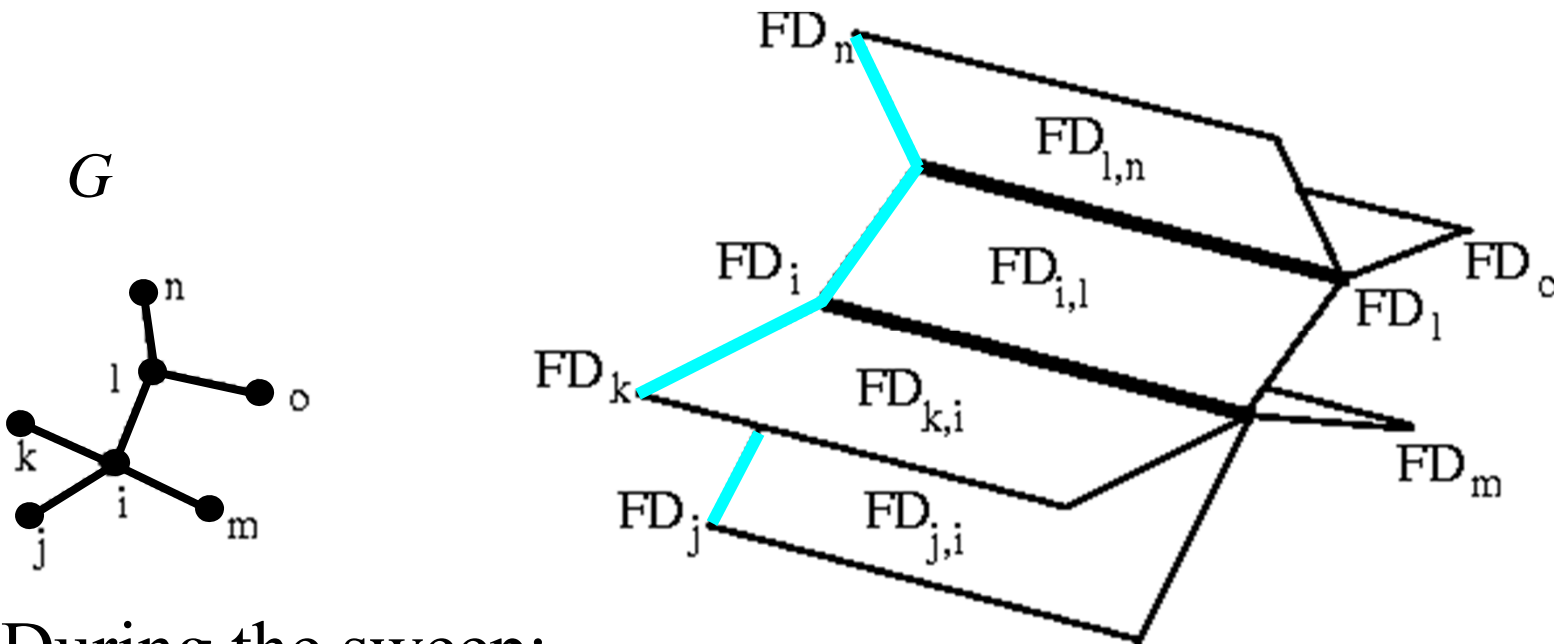
- Sweep all $FD_{i,j}$ with a **sweep line** from left to right

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Compute Reachable Points



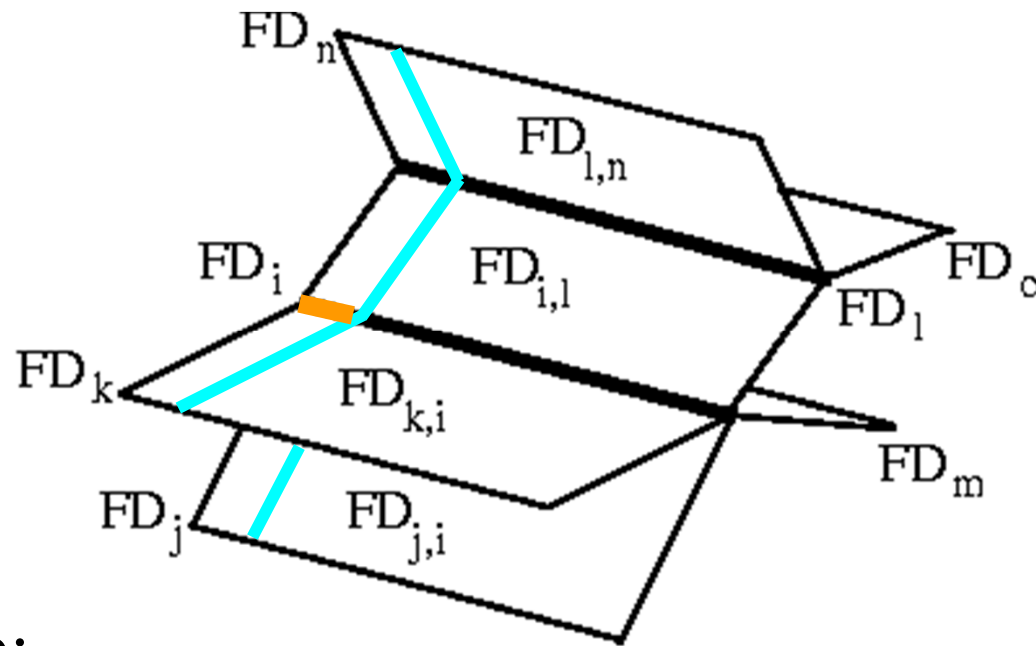
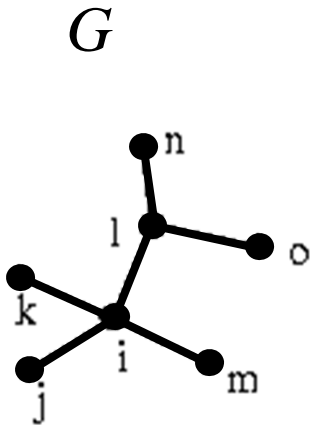
- During the sweep:
Compute **points** on the free space surface, to the left of the sweep line, which are reachable by a monotone path from a lower left corner.

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Compute Reachable Points



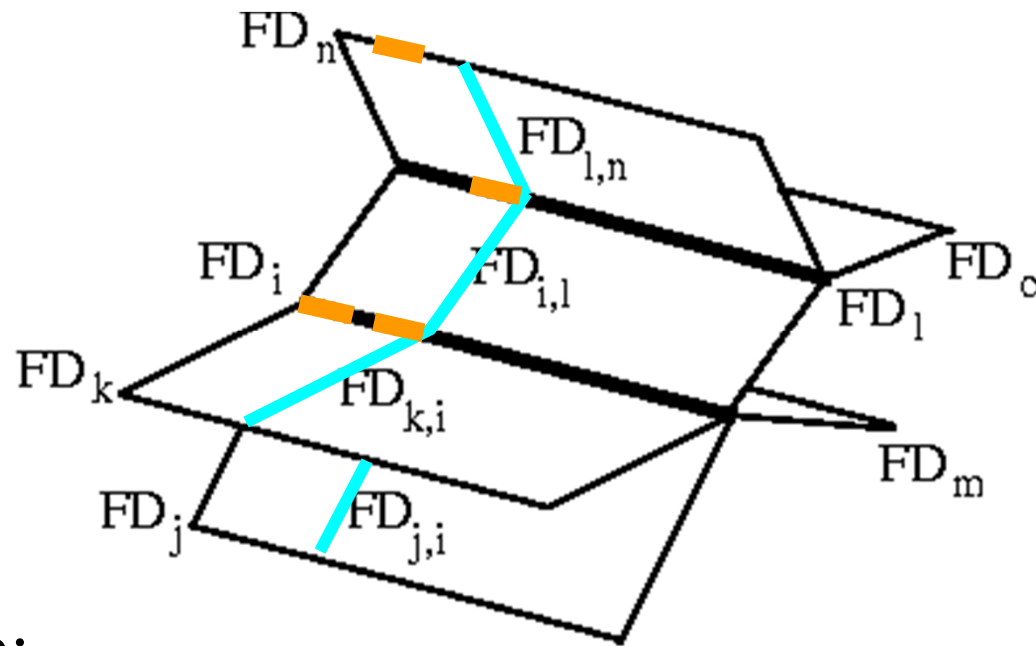
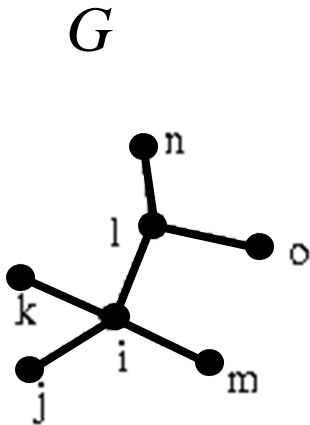
- During the sweep:
Compute **points** on the free space surface, to the left of the sweep line, which are reachable by a monotone path from a lower left corner.

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Compute Reachable Points



- During the sweep:

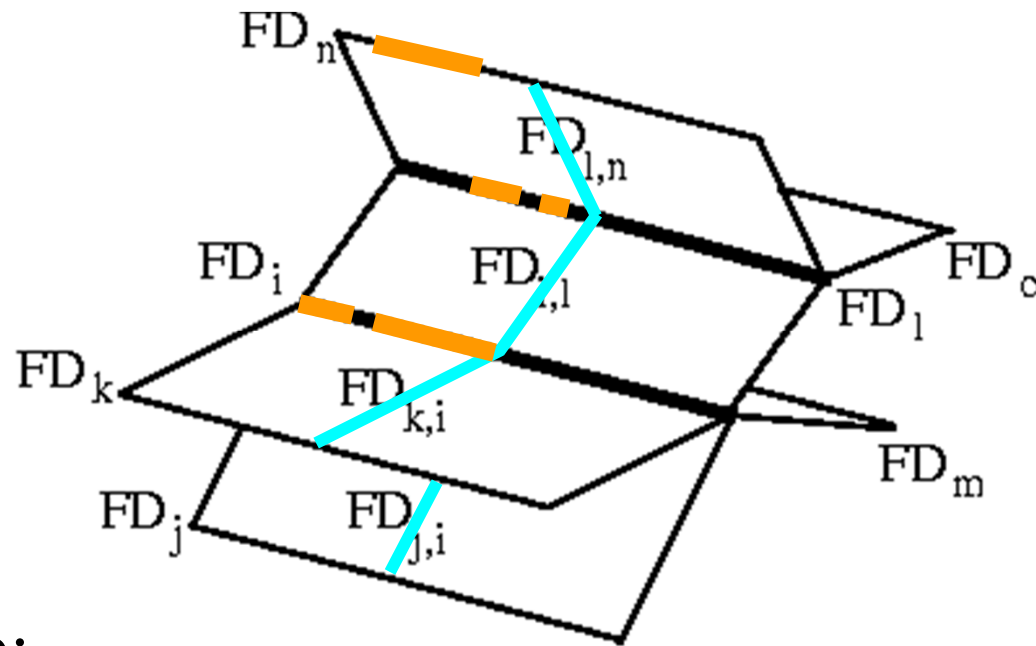
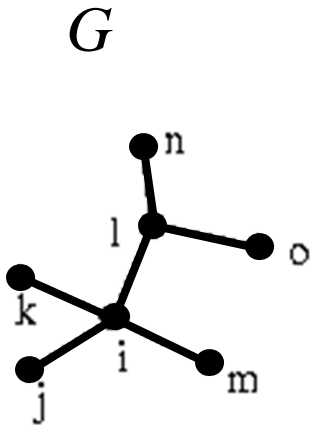
Compute **points** on the free space surface, to the left of the sweep line, which are reachable by a monotone path from a lower left corner.

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Compute Reachable Points



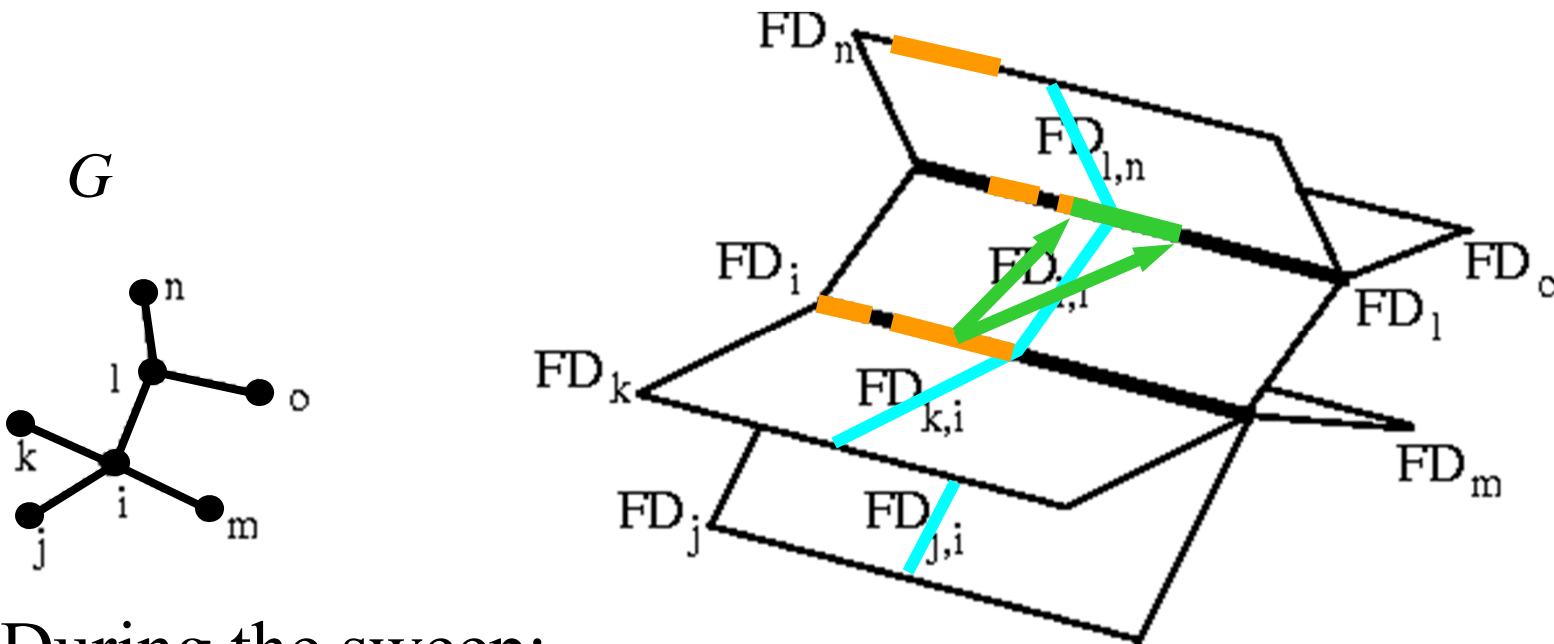
- During the sweep:
Compute **points** on the free space surface, to the left of the sweep line, which are reachable by a monotone path from a lower left corner.

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



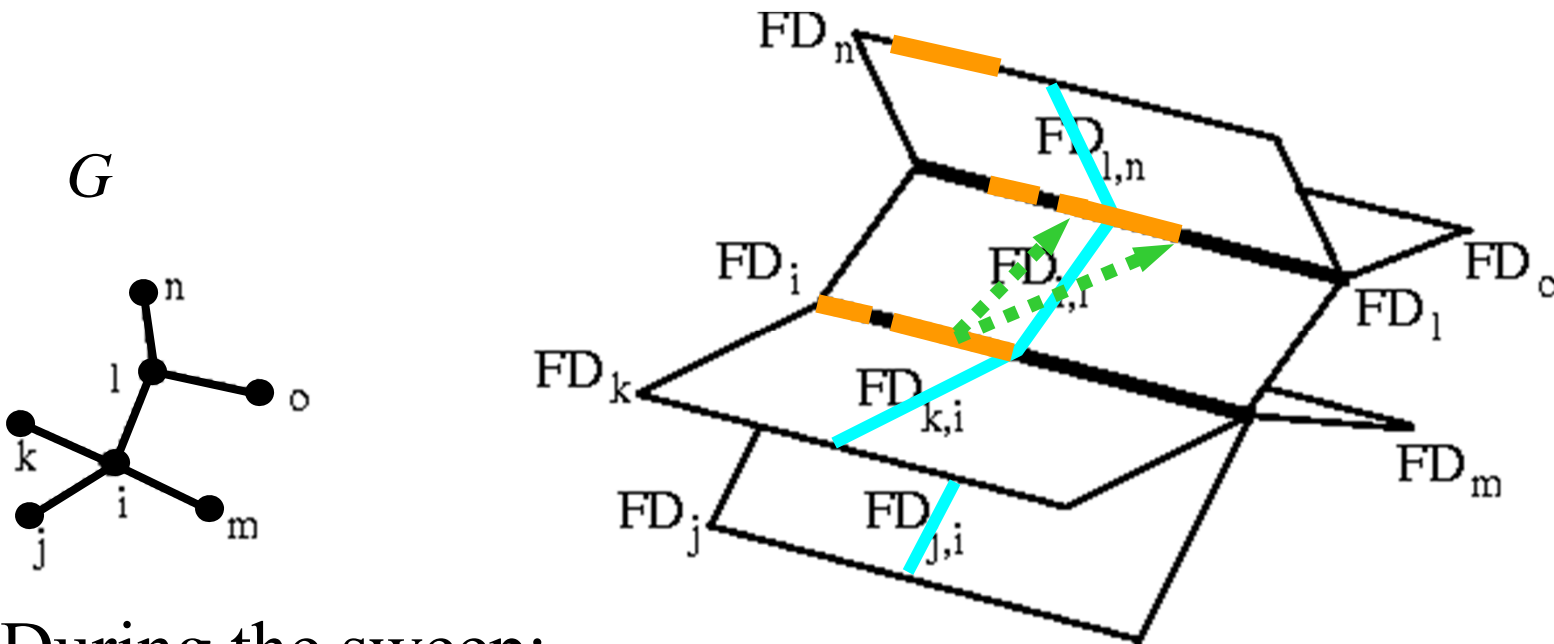
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



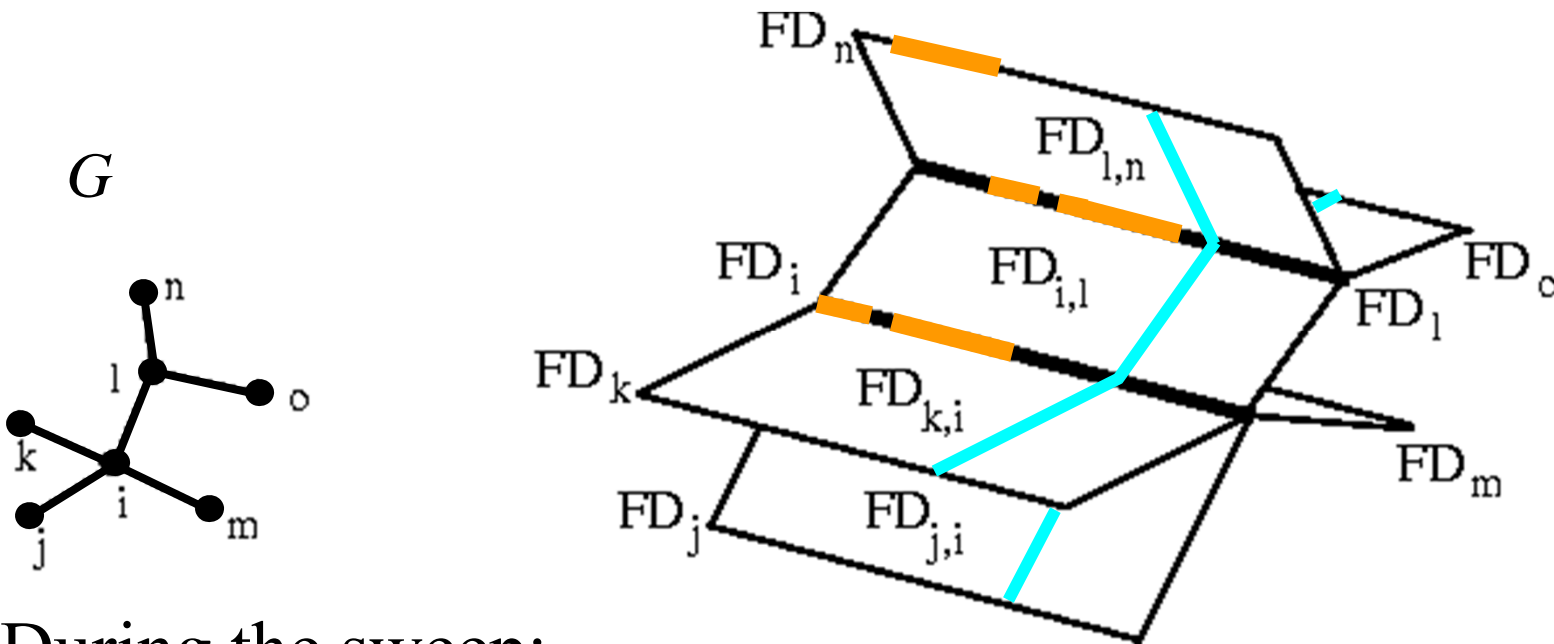
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



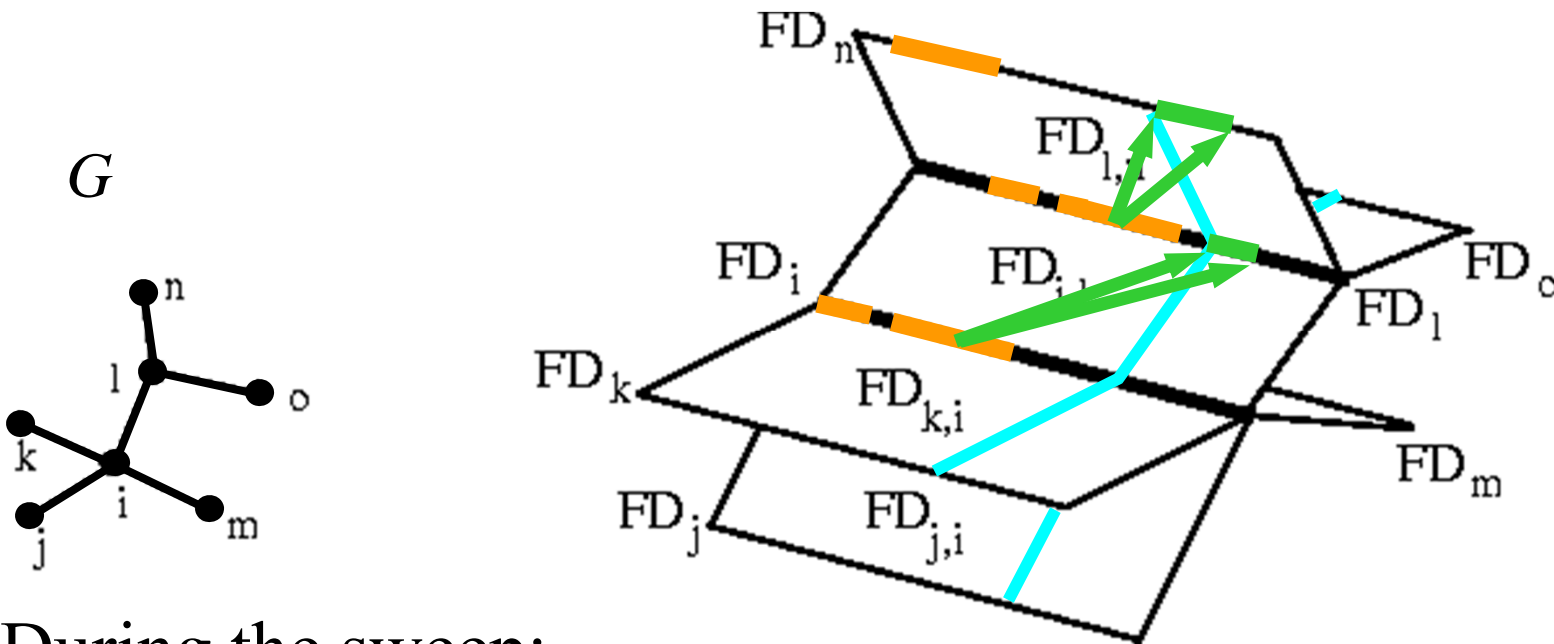
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



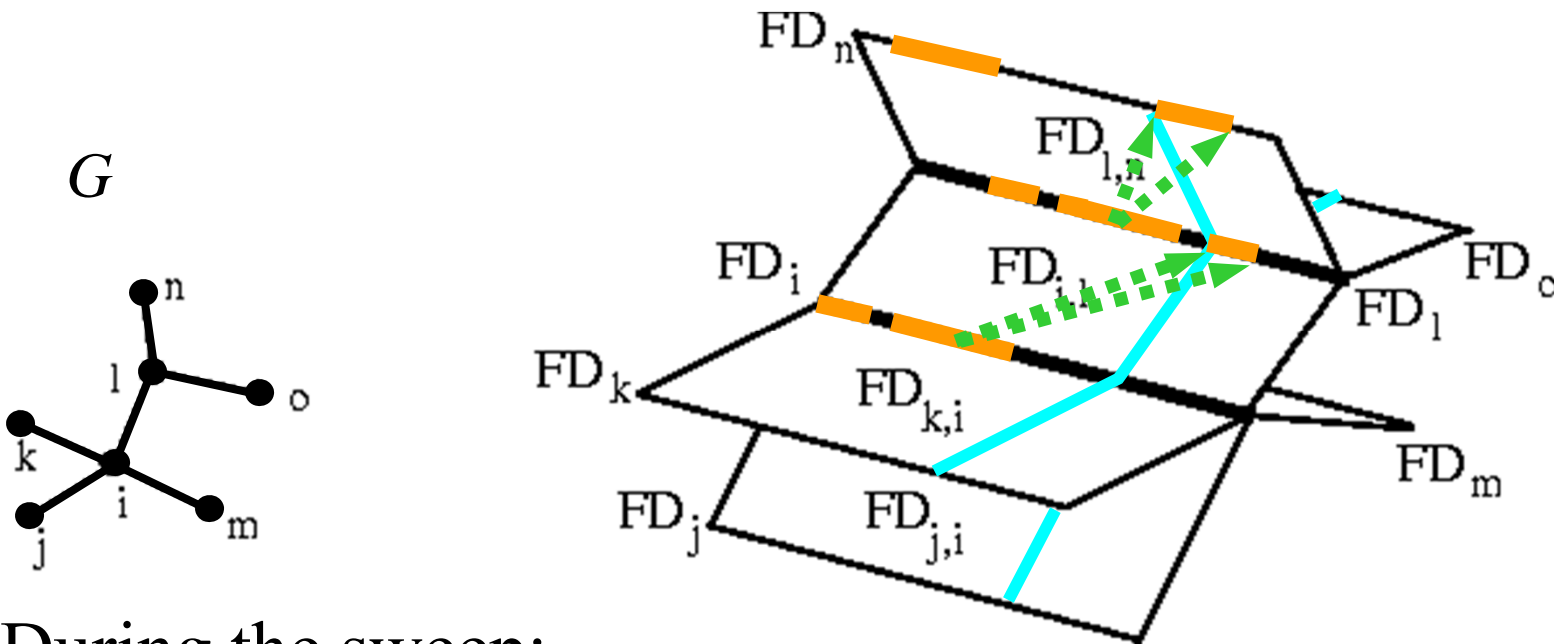
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



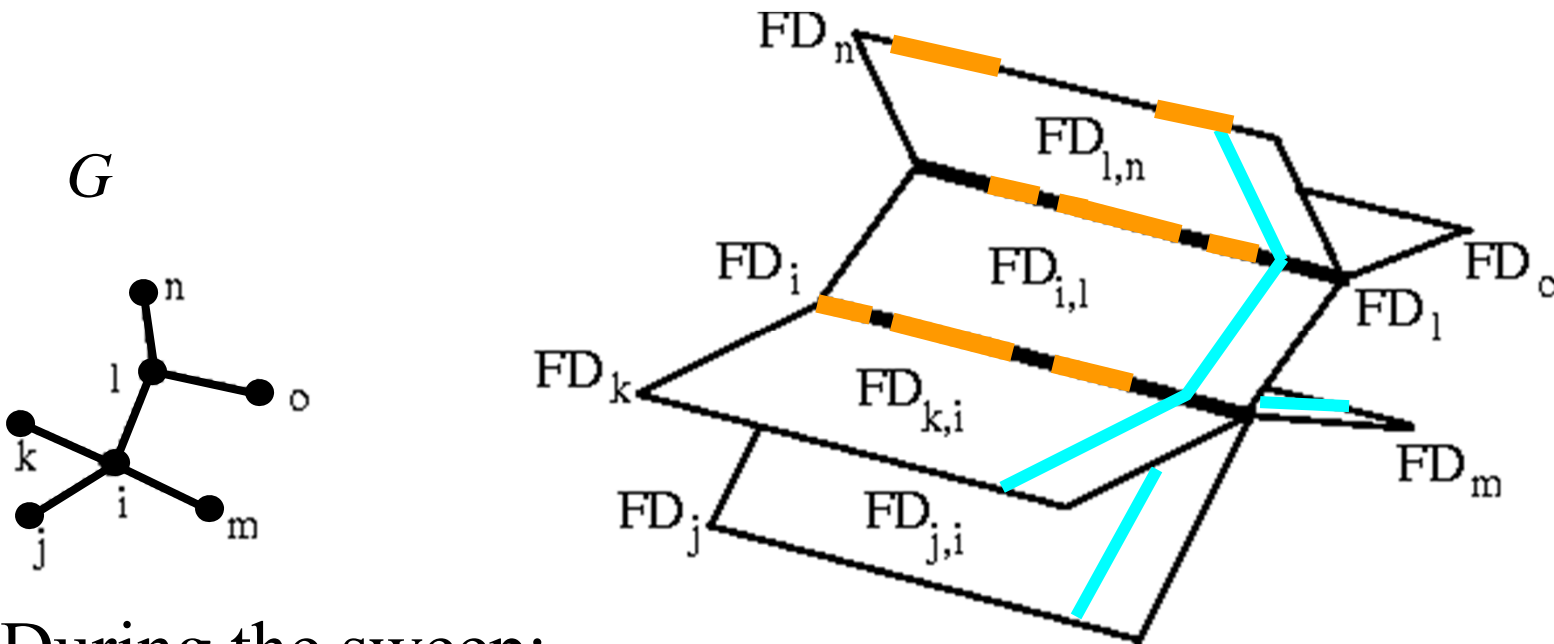
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



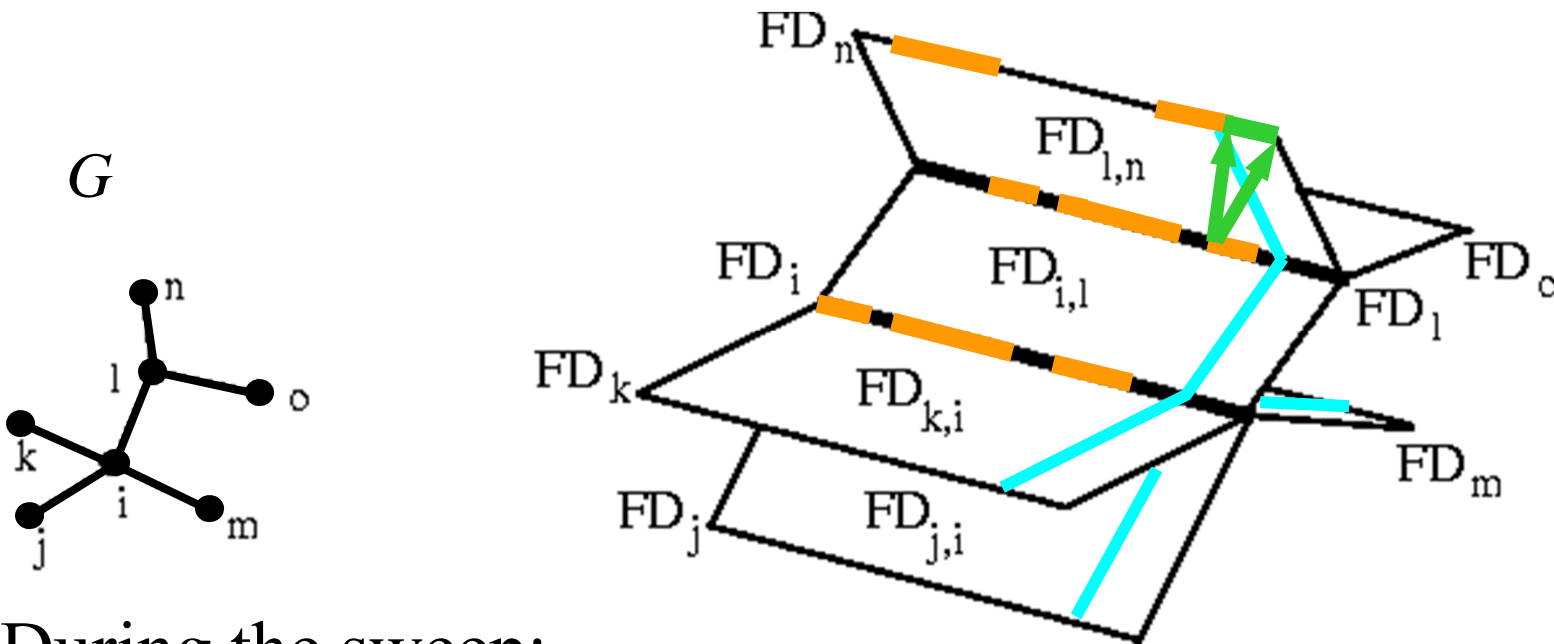
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



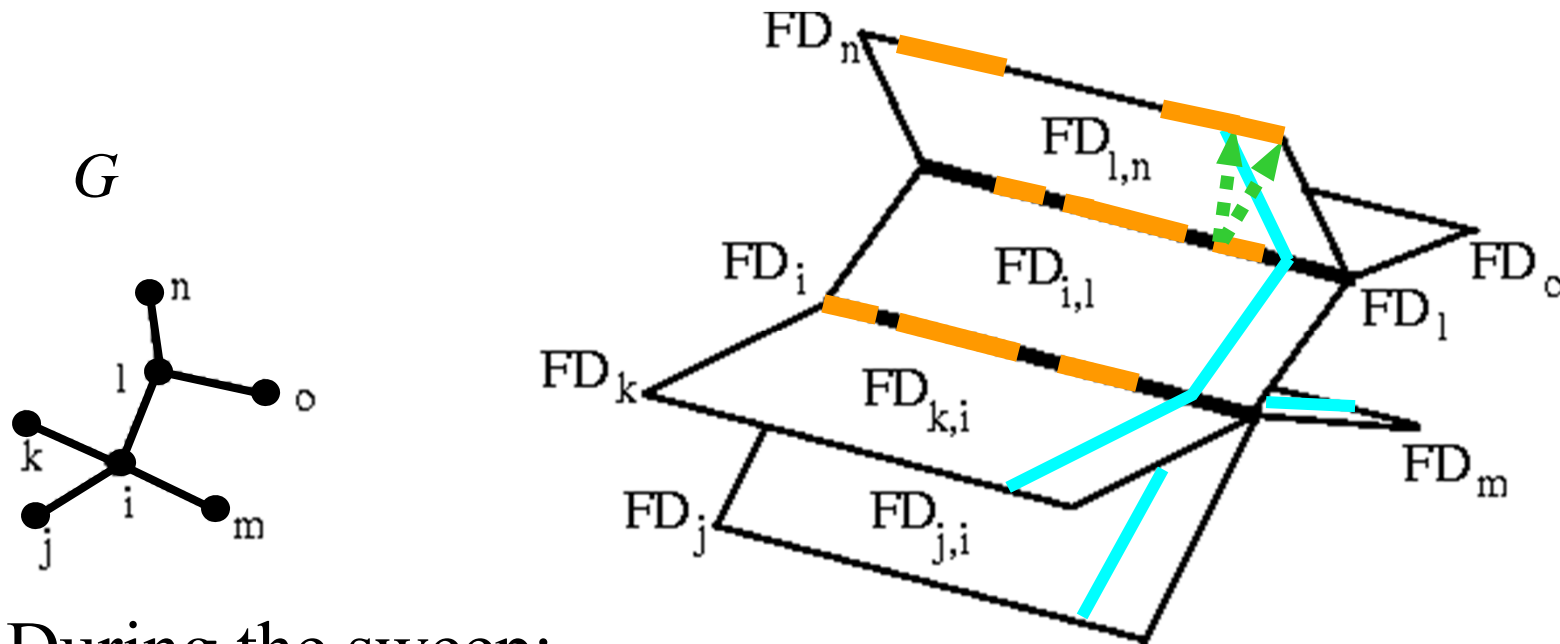
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



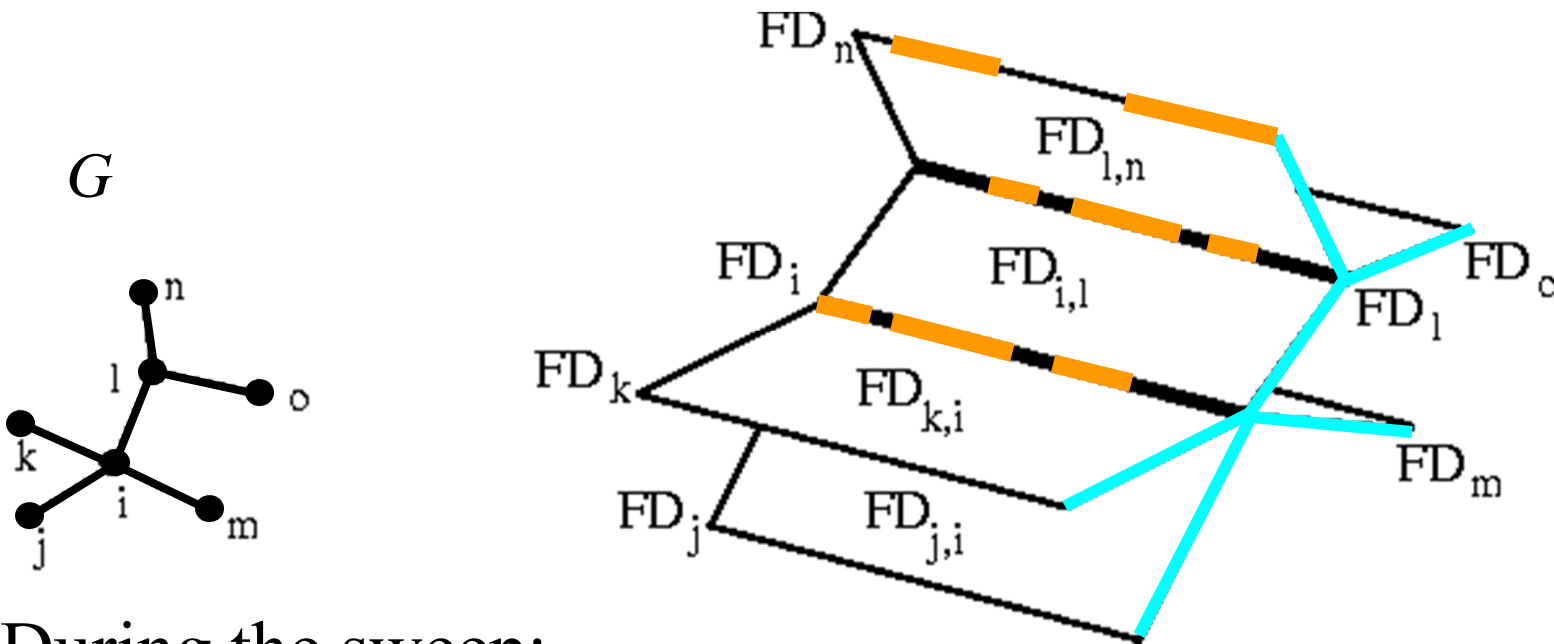
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Update Reachable Points



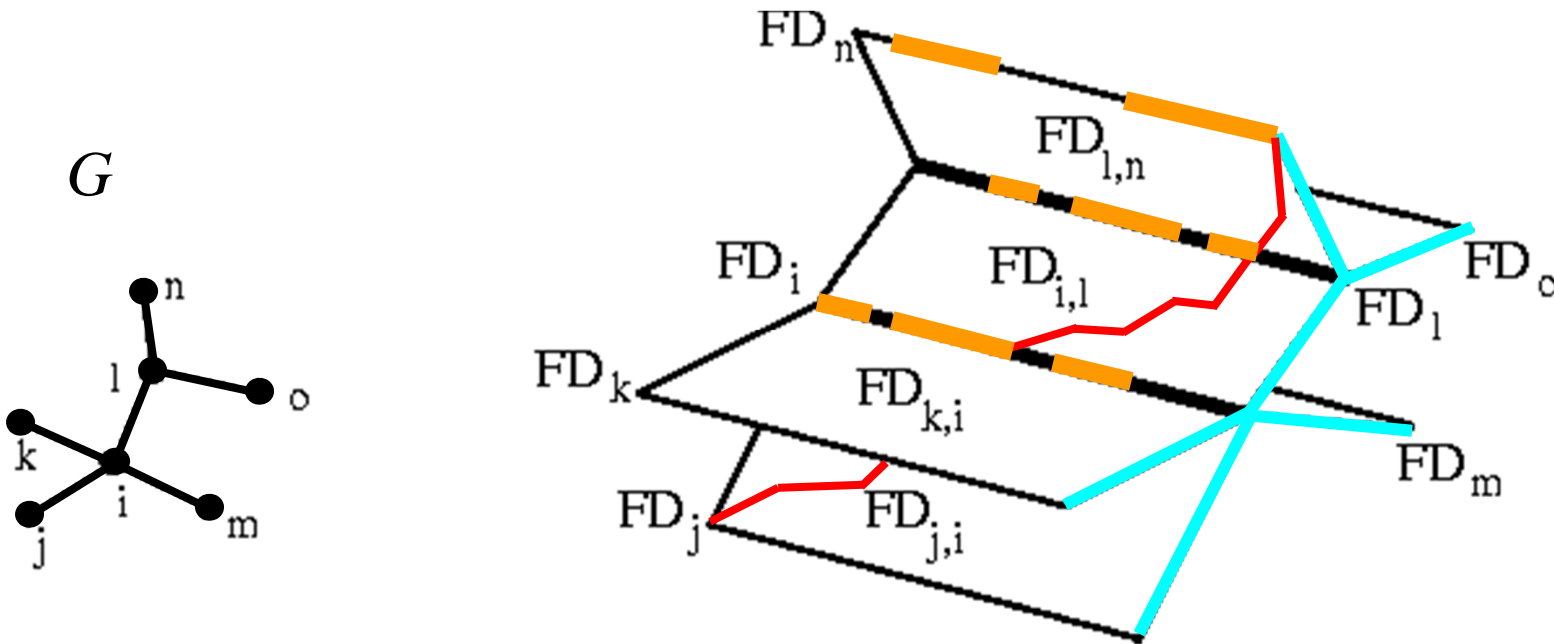
- During the sweep:
 - Update reachable points Dijkstra-style
 - Use a data structure which supports reachability queries in the free space surface

[AERW03] H. Alt, A. Efrat, G. Rote, C. Wenk, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, C. Wenk, On Map-Matching Vehicle Tracking Data, *VLDB* 853-864, 2005.

[WSP06] C. Wenk, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., *SSDBM*: 379-388, 2006.

Backtracking



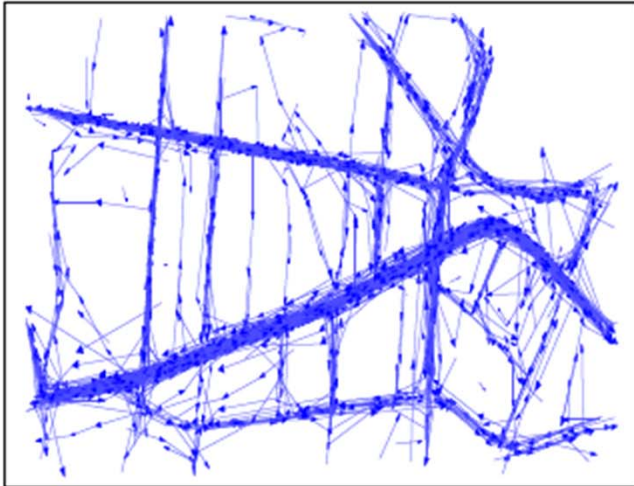
- After the sweep:
 - Construct a **monotone path via** backtracking
- Runtime:
 - $O(mn \log(mn))$ time for decision problem; optimize using parametric search

[AERW03] H. Alt, A. Efrat, G. Rote, **C. Wenk**, Matching Planar Maps, *J. of Algorithms* 49: 262-283, 2003.

[BPSW05] S. Brakatsoulas, D. Pfoser, R. Salas, **C. Wenk**, On Map-Matching Vehicle Tracking Data, VLDB 853-864, 2005.

[WSP06] **C. Wenk**, R. Salas, D. Pfoser, Addressing the Need for Map-Matching Speed..., SSDBM: 379-388, 2006.

Map Construction



Map Construction [AW12]

We model the original map and the reconstructed map as embedded undirected graphs in the plane.

We model error associated with each trajectory by a precision parameter ε .

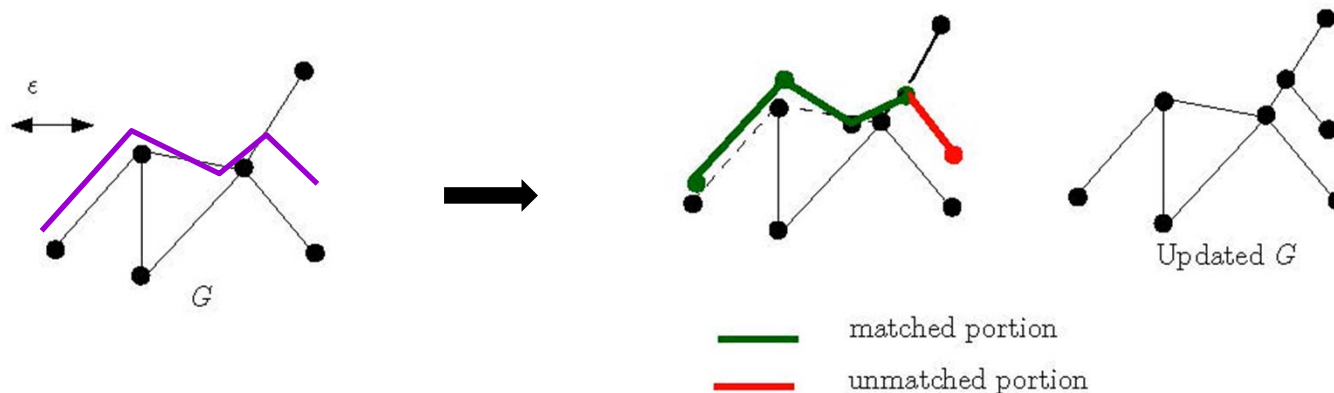
1. We assume each input curve is within Frechet distance $\varepsilon/2$ of a street-path in the original map.
2. (We assume all input curves sample acyclic paths.)
3. Two additional assumptions on original map help us to provide quality guarantees.

Map Construction [AW12]

Incrementally add one trajectory after another.

For each trajectory:

1. Use partial Fréchet distance to identify new and existing portions by combining mapmatching with partial Fréchet distance:



2. Use min-link curve simplification algorithm to reconcile existing portions

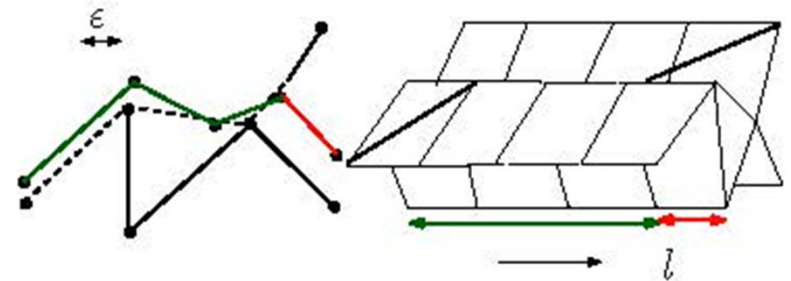
Map Construction [AW12]

Incrementally add one trajectory after another.

For each trajectory:

1. Use partial Fréchet distance to identify new and existing portions by combining mapmatching with partial Fréchet distance:

- Compute free space surface
- Find path that maximizes matched portion on the curve.



⇒ Project free space onto curve:

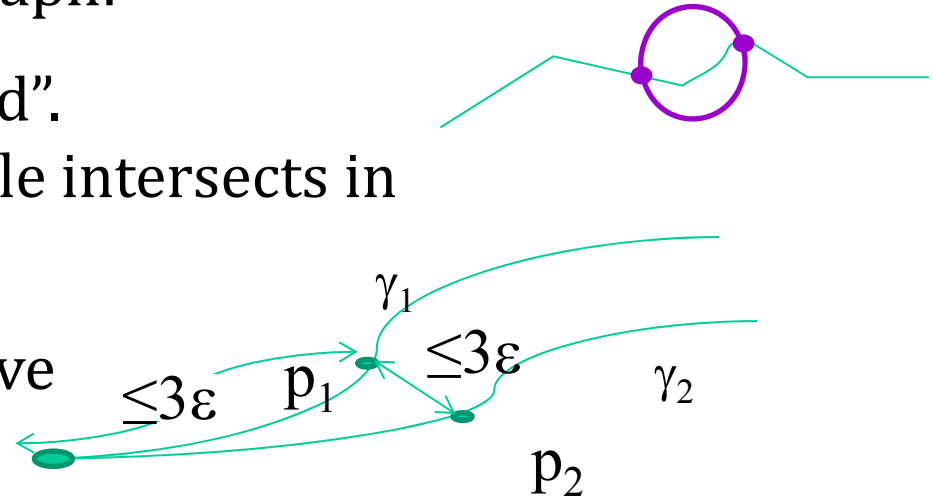
white interval = matched portion, black interval = unmatched portion

2. Use min-link curve simplification algorithm to reconcile existing portions

Assumptions

Assumptions on unknown graph:

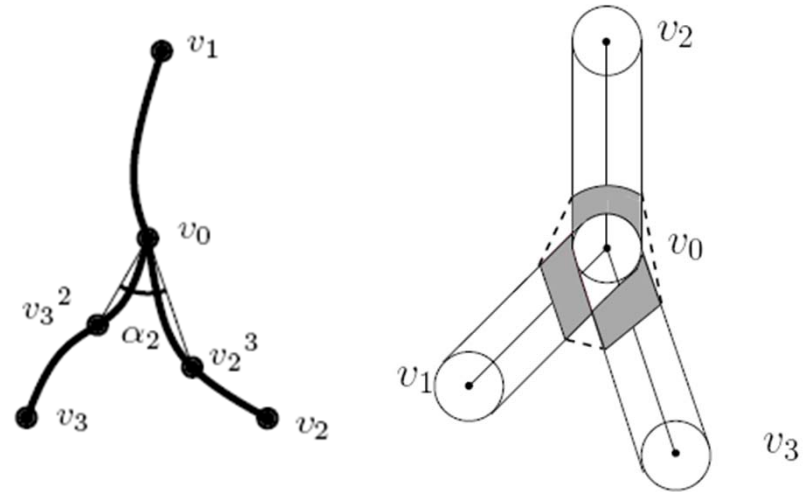
1. Road fragments are "good".
"good": Every small circle intersects in just two points
2. Close fragments must have an intersection point



\Rightarrow Projection approach is justified, because free space has special structure. Trajectory can only sample one good section in original network.

Give quality guarantees

- **Good regions:** We prove the quality guarantee that there is a 1-to-1 correspondence with bounded description complexity between well-separable good portions of original network and reconstructed graph.
- **Bad regions:** We give the first description and analysis of vertex regions.



⇒ It is relatively easy to handle well-sampled clean data.
Deal with noisy data that is not well-sampled and give quality guarantees.

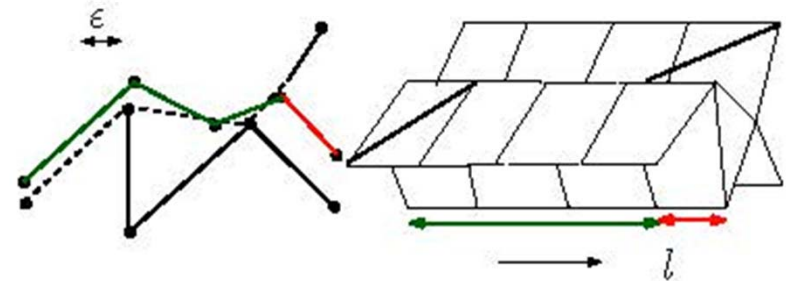
Map Construction [AW12]

Incrementally add one trajectory after another.

For each trajectory:

1. Use partial Fréchet distance to identify new and existing portions by combining mapmatching with partial Fréchet distance:

- Compute free space surface
- Find path that maximizes matched portion on the curve.



⇒ Project free space onto curve:

white interval = matched portion, black interval = unmatched portion

2. Use min-link curve simplification algorithm to reconcile existing portions

Curve Simplification

- **Arbitrary vertices:**

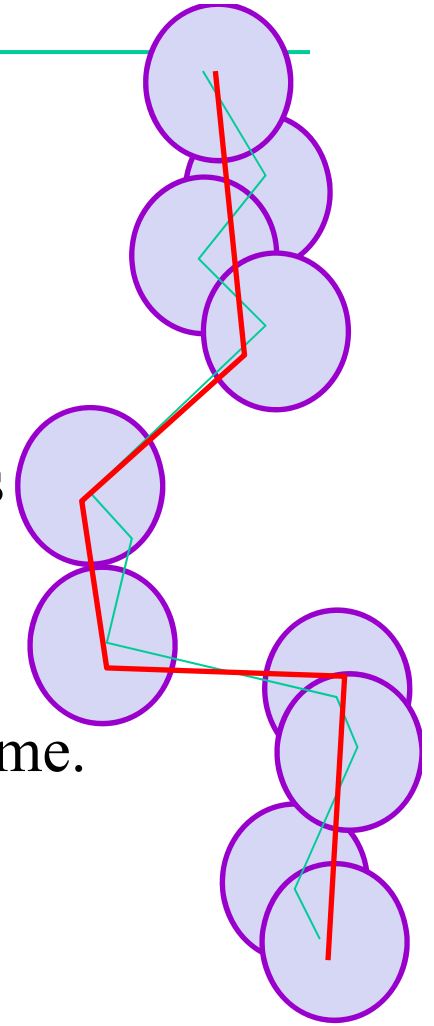
- **min-#:** Compute a minimum link path that stabs ε -neighborhoods around vertices of f in the correct order. $O(n)$ time.

- **min- ε :** Binary search on critical values. $O(n^2)$ time.

- **Subsequence of original vertices:**

- Douglas-Peucker algorithm, $O(n \log n)$ time.

- **min-#:** Add shortcut-edges between vertices if shortcut is within error ε . Then find a path with min-# edges. $O(n^2)$ time, and $O(n^{4/3+\delta})$ for x-monotone chains.



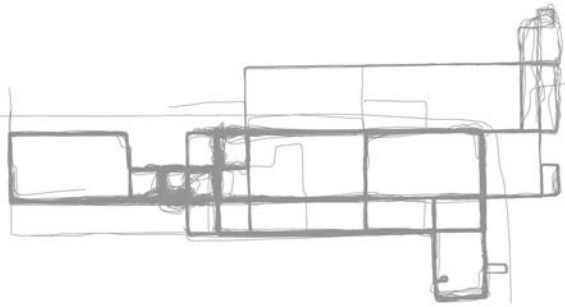
[AG99] H. Alt, L. Guibas, Discrete Geometric Shapes..., Handbook of Computational Geometry: 121-153, 1999.

[GHMS93] L. Guibas, J. Hershberger, J. Mitchell, J. Snoeyink, Approximating ..., minimum link paths, IJCGA 3(4): 383-415, 1993

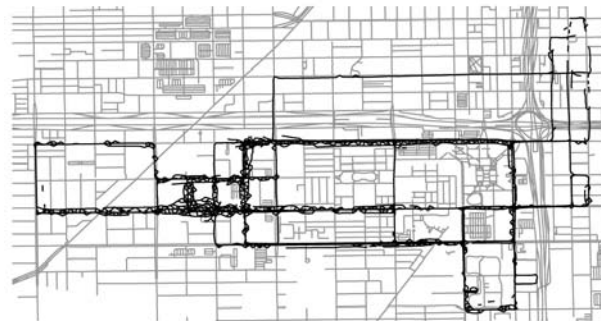
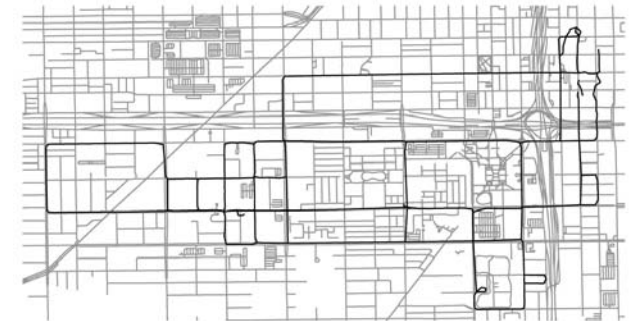
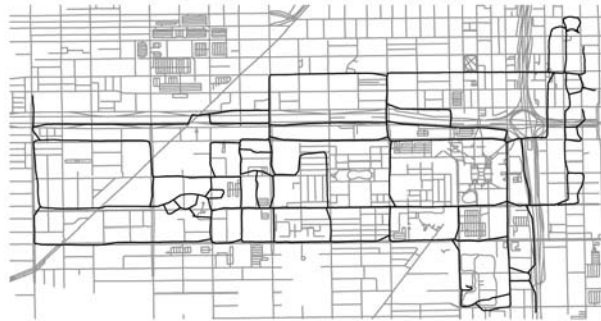
[II88] H. Imai, M. Iri, Polygonal approximations of a curve..., in Computational Morphology, Toussaint (eds): 71-86, 1988

Different Reconstructions

GPS Trajectory Data



Reconstructed Maps



Which is the Better Roadmap ?



Which is the Better Roadmap ?



⇒ Map Comparison

Distance Measures for Map Comparison

- [BE12b]: Graph sampling-based distance measure in local neighborhood.
Maximize number of marbles and holes that match 1-to-1.
- [KP12]: Compare shortest paths in both maps, with nearby start and end positions. Ensures similar connectivity/routing properties.
- [BE12], [AKPW14]: Overview / benchmark papers

[AFHW14]: Considers maps as sets of paths, and compares path sets.

[AFW13]: Compares local topology of graphs using persistent homology

[AFHW14] M. Ahmed, B. Fasy, K. Hickmann, **C. Wenk**, Path-based distance for street map comparison, arXiv:1309.6131, 2014.
[AFW13] M. Ahmed, B. Fasy, **C. Wenk**, Local homology based distance between maps, submitted to SoCG, 2013.

[KP12] S. Karagiorgou, D. Pfoser, On vehicle-tracking data-based road network generation, 20th ACM SIGSPATIAL: 89-98, 2012.
[BE12] J. Biagioni, J. Eriksson, Map inference in the face of noise and disparity, 20th ACM SIGSPATIAL: 79-88, 2012.
[BE12b] J. Biagioni, J. Eriksson, Inferring road maps from global... TRR: J. of the Transportation Research Board 2291, 61-71, 2012.
[AKPW14] M. Ahmed, S. Karagiorgou, D. Pfoser, **C. Wenk**, A Comparison and Evaluation of ..., submitted to Geoinformatica, 2014.

[AFHW14] Path-Based Distance

- Directed Hausdorff distance on path-sets:

$$\vec{d}_{G,H}(\pi_G, \pi_H) = \max_{p_G \in \pi_G} \min_{p_H \in \pi_H} \delta_F(p_G, p_H)$$

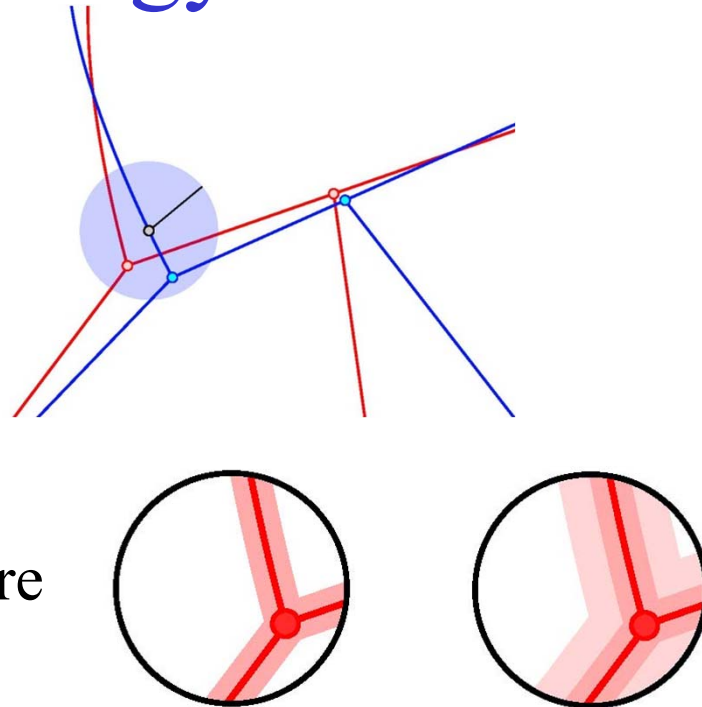
- π_G path-set in G , and π_H path-set in H
- We prove that using the set of paths of link-length three approximates the overall distance, if vertices in G are well-separated and have degree $\neq 3$.
- Asymmetry of distance definition is desirable, if G is a reconstructed map and H a ground-truth map.

[AFHW14] M. Ahmed, B. Fasy, K. Hickmann, **C. Wenk**, Path-based distance for street map comparison, arXiv:1309.6131, 2014.
[AFW13] M. Ahmed, B. Fasy, **C. Wenk**, Local homology based distance between maps, submitted to SoCG, 2013.

[KP12] S. Karagiorgou, D. Pfoser, On vehicle-tracking data-based road network generation, 20th ACM SIGSPATIAL: 89-98, 2012.
[BE12] J. Biagioni, J. Eriksson, Map inference in the face of noise and disparity, 20th ACM SIGSPATIAL: 79-88, 2012.
[BE12b] J. Biagioni, J. Eriksson, Inferring road maps from global... TRR: J. of the Transportation Research Board 2291, 61-71, 2012.
[AKPW14] M. Ahmed, S. Karagiorgou, D. Pfoser, **C. Wenk**, A Comparison and Evaluation of ..., submitted to Geoinformatica, 2014.

[AFW13] Local Homology Based Distance

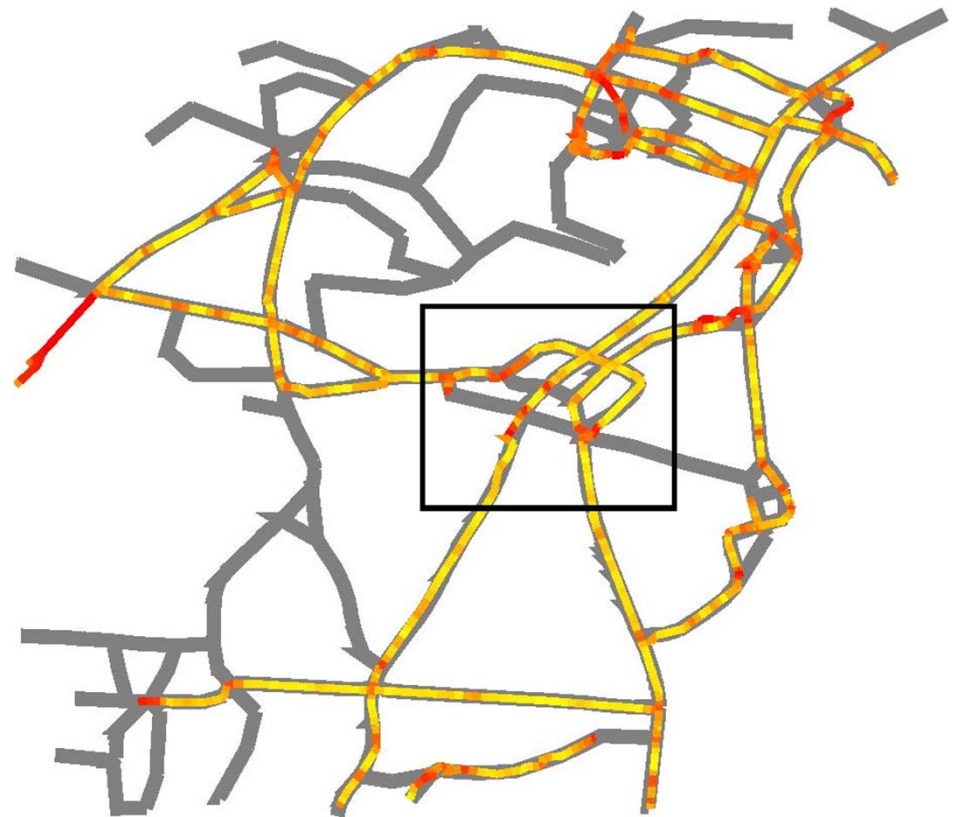
- Consider a common local neighborhood of both maps.
- Consider the cycles of each graph inside this neighborhood.
- Now thicken each graph and track changes in the cycle structure using persistent homology

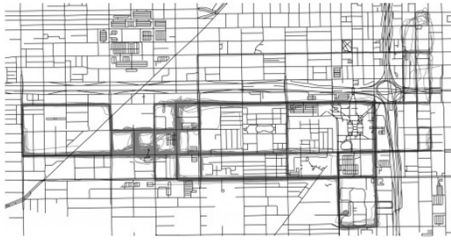


- ⇒ Use distance between persistence diagrams to compare changing local cycle structure
- ⇒ Local “signature” that captures local topological similarity of graphs

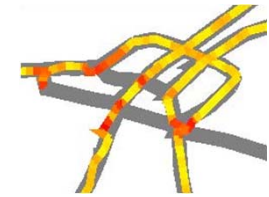
[AFW13] Local Homology Based Distance

- Compared two reconstructed maps.
- Local signature captures different topology (missing intersections) well





Conclusion



- Map construction and map comparison are recent data-driven problems
- Related to geometric reconstruction, trajectory clustering, shape comparison
- There is a lot of potential for theoretical modeling and algorithms that provide quality guarantees
- Open problems / future work:
 - Map updates
 - More complicated/realistic input and noise models for trajectories
 - More complicated/realistic output models for the maps (vertex regions; directed graphs, with turn information, road categories, etc.)

